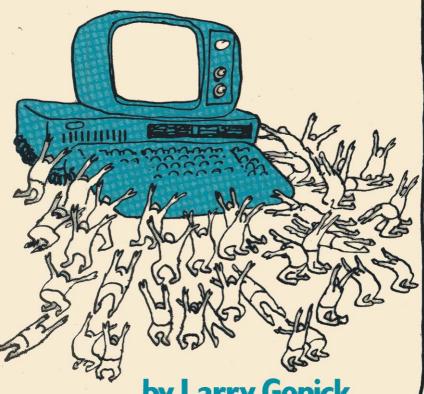
THE CARTOON GUIDE TO

COMPUTER SCIENCE



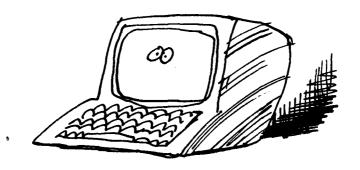
by Larry Gonick

Also available by Larry Gonick, with Mark Wheelis THE CARTOON GUIDE TO GENETICS

THE CARTOON GUIDE TO

COMPUTER SCIENCE

Larry Gonick





BARNES & NOBLE BOOKS

A DIVISION OF HARPER & ROW, PUBLISHERS

New York, Cambridge, Philadelphia

San Francisco, London, Mexico City

São Paulo, Sydney

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FIRST EDITION

Library of Congress Cataloging in Publication Data

Gonick, Larry.

The cartoon guide to computer science.

(College outline series : CO/417)

Includes index.

Summary: An introduction to computer science in cartoon format.

 Computers—Caricatures and cartoons.
 Electronic data processing— Caricatures and cartoons.
 Computers—Cartoons and comics.
 Data processing—Cartoons and comics.
 Title.
 Series.

QA76.G593 1983 001.64 82-48251

ISBN 0-06-460417-9 (pbk.)

83 84 85 86 87 10 9 8 7 6 5 4 3 2 1

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LARRY GONIUM, THE OVEREDUCATED CARTOONIST, HOLDS TWO DEGRESS IN MATHEMATICS FROM HARVARD. HE HAS WORKED AS A FORTRAN PROGRAMMER, AND SOME OF HIS BEST FRENDS ARE IN THE CONPUTER BUSINESS. HE LIVES IN SAN FRANCISCO WITH HIS NIFE AND DAUGHTER, WHO WOULD LIKE TO FIND SOME CARTOON PROCESSING SOFTWARE TO IMPROVE HIS PRODUCTIVITY.



COMPUTER SCIENCE

LEARNING HAS NEVER BEEN SO EASY OR SO MUCH FUN

Here are the elements of computer science illustrated, simplified, and humor-coated so that you understand them at once. Use this book to lighten up that serious course you are taking or to penetrate the fog of that equally serious textbook you are trying to follow. Read it to gain both an overview and an inner view of that computer you are learning to use. Or if you feel the computer revolution is passing you by, let it give you a point of entry. It won't make a programmer out of you, but it will put you well on the way to computer literacy.

In these pages you'll meet Charles Babbage and his analytical engine, which was never built, and Ada Augusta, Lady Lovelace, who programmed it nevertheless. You'll also meet George Boole, whose algebra underlies the design of circuitry. You'll learn about binary numbers, computer components and architecture, software, programming languages from machine language to BASIC, and special computer applications—cryptography, artificial intelligence, and others you may not have heard of.

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Cover design by Larry Gonick

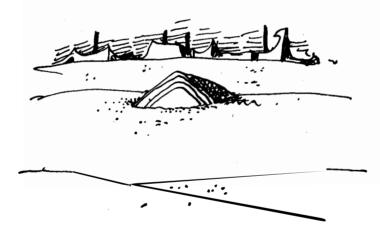
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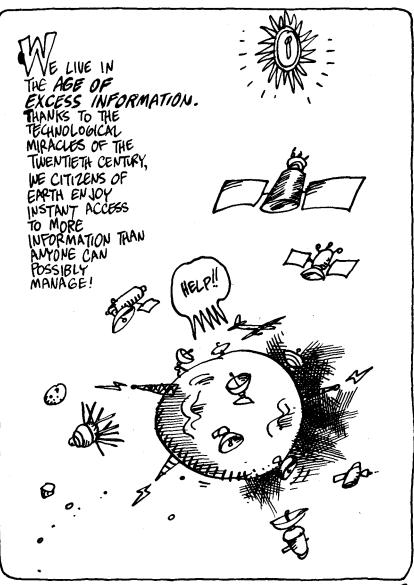
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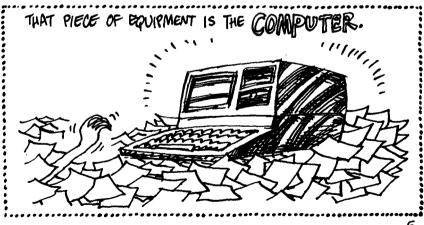
PART I THE AGES OF INFORMATION

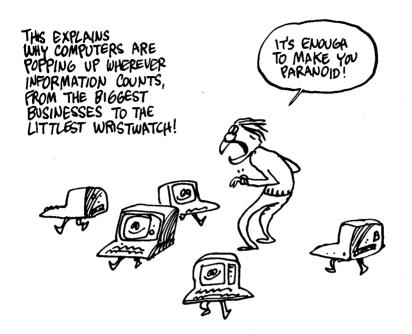














What is information?

IN THE EVERYDAY
SENSE OF THE WORD,
"WFORMATION" MEANS
FACTS: THE SORT OF
STUFF THAT FILLS
NON FICTION BOOKS, AND
CAN ONLY BE EXPRESSED
IN WORDS.



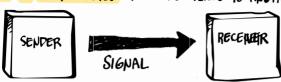
IN THE WORLD OF COMPUTERS, HOWEVER, THE TERM HAS A MUCH BROADER MEANING.



THE MODERN
DEFINITION COMES
FROM CLAUDE
SHANNON, A
BELL LABS
ENGINEER,
AMATEUR UNICYCUST,
AND FOUNDER OF
THE SCIENCE OF
INFORMATION
THEORY.

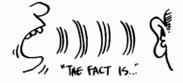
SHANNON ALSO BUILT AN ELECTRIC "MOUSE" THAT COULD BE PROGRAMMED TO RUN MAZES!

ACCORDING TO SHANYON, INFORMATION IS PRESENT WHENEVER A SIGNAL IS TRANSMITTED FROM ONE PLACE TO ANOTHER:



IT DOESN'T MATTER WHAT KIND OF SIGNAL IT IS. FOR EXAMPLE:

THE SIGNAL MAY BE IN THE FORM OF WORDS, THE MOST FAMILIAR KIND OF WFORMATION...

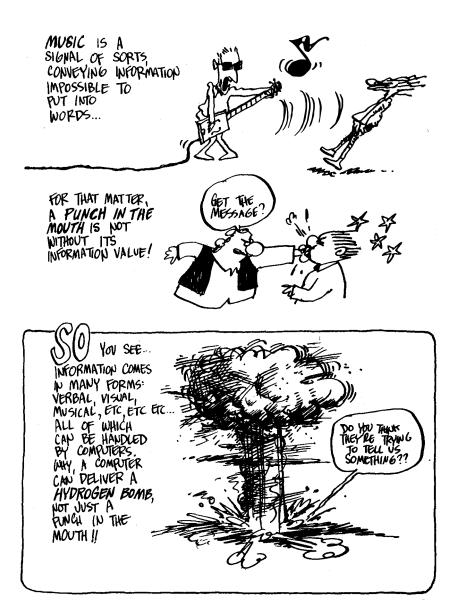


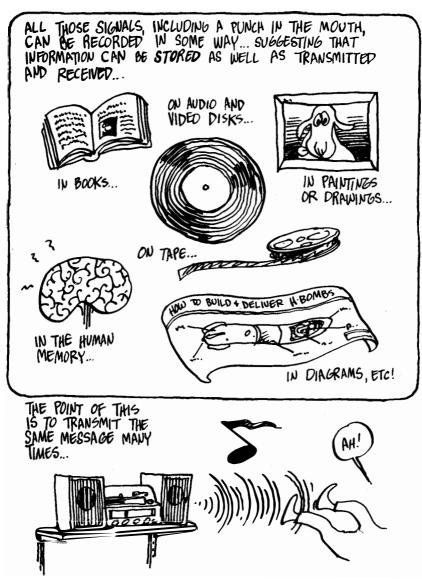
...BUT A PICTURE
ALSO SENDS A
SIGNAL, IN THE
FORM OF LIGHT
WAVES, TO OUR EYES.
IT LOOKS AS IF PICTURES
CONNEY INFORMATION!



FURTHERMORE, OUR EYE SAYDS A PATTERN OF ELECTRIC IMPULSES UP THE OPTIC NERVE TO THE BRAIN. THAT SIGNAL CARRIES INFORMATION, TOO!

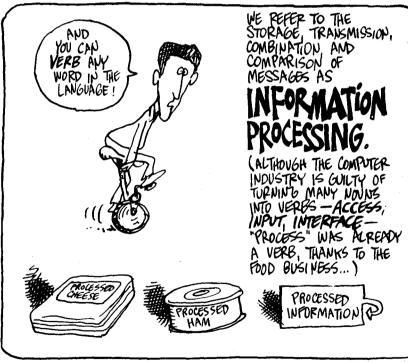






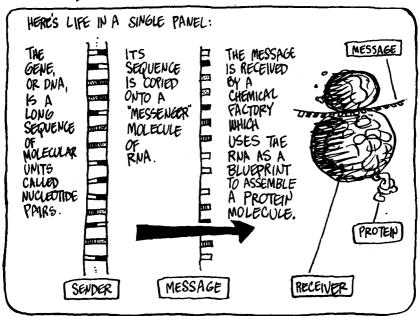
AND OF COURSE, ITEMS OF INFORMATION CAN BE COMBINED IN VARIOUS WAYS.





TO APPRECIATE THE POWER OF INFORMATION, CONSIDER ANOTHER EVERYDAY EXAMPLE:

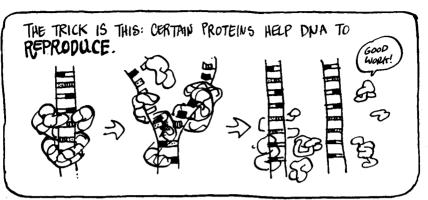
LIFE ITSELF.



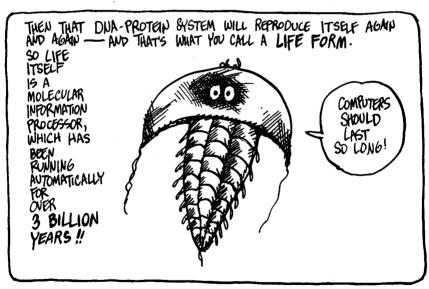


IN OTHER WORDS, THE PROTEIN IS BUILT ACCORDING TO INFORMATION STORED IN THE GENE.





WHAT HAPPENS THEN? IF DNA ENCODES PROTEINS THAT HELP DNA TO REPRODUCE, THEN MORE OF THOSE PROTEINS WILL BE BUILT, MORE DNA WILL BE COPIED...ETC! MOREOVER, IF THE DNA ENCODES OTHER PROTEINS WHICH PROTECT IT IN VARIOUS WAYS, AND OTHERS TO ATTACK AND DESTROY RIVAL DNA AND PROTEINS...



The Evolution of the Computer

IT MAY BE GOING TOO FAR TO SAY THAT COMPUTERS HAVE BEEN EVOLVING FROM THE BEGINNING...

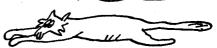


BUT FROM EARLY TIMES, LIFE FORMS HAVE BEEN INCREASING THEIR INFORMATION-PROCESSING ABILITIES. EVEN AN AMORBA RECEIVES CHEMICAL SIGNALS TELLING IT WHERE THE FOOD IS!

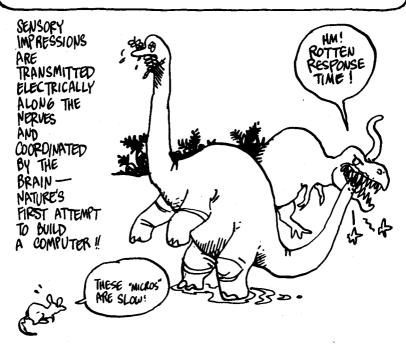




BUT THEN WE CALL IT PURR-CEIVING!

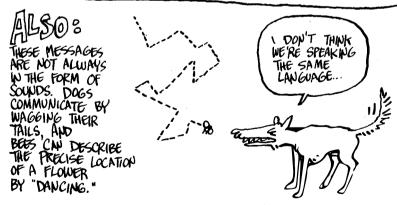


THE BYES PERCEIVE A RANGE OF ELECTROMAGNETIC RAYS; THE EARS RESPOND TO PRESSURE IN THE AIR; THE NOSE REACTS TO VARIOUS MOLECULES; SO DO THE TASTE BUDS; AND THE SENSE OF TOUCH IS A WAY OF RECEIVING A PUNCH IN THE MOUTH!

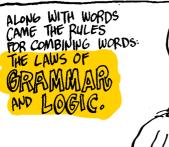


BESIDES TRANSMITTING INFORMATION WITHIN THEIR OWN BODIES, ANIMALS ALSO SENT MESSAGES TO EACH OTHER: KKK NOTE AGAIN: THESE DO NOT WHAT DO NECESSARILY YOU MEAN, WHAT CONVEY INFORMATION DOES "GRONK" MEAN? "GRONK THAT CAN BE EXPRESSED IN MEANS "GROUK!" WORDS! -YAWYUA WHAT DOES

"MEAN" NEAN?







IF YOU COME OUT AND APOLOGIZE, THEN WE WILL NOT FLAY YOU ALIVE, UNLESS WE CHANGE OUR MINDS...



IN TIME, YOWEVER, IT APPEARED THORE WAS A SPECIAL TYPE OF WORD WITH ITS OWN SPECIAL RULES... NAMBLY—



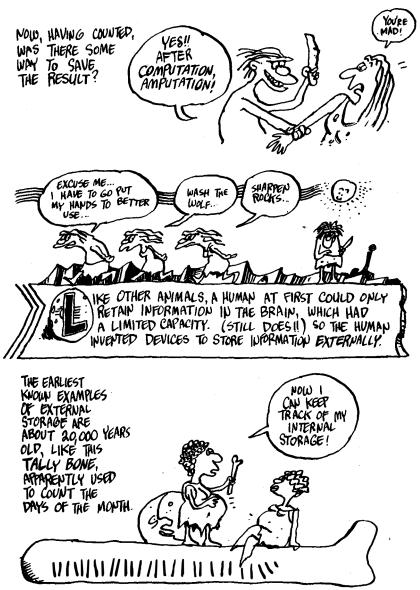
WAIT ONE MINUTE... LET GUESS...



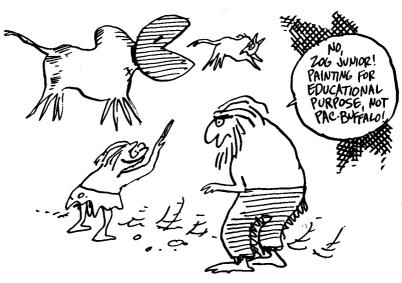
NUMBERS ARE PRECISE... RELIABLE...
YOU CAN ADD, SUBTRACT AND
MULTIPLY NUMBERS... "ONE PLUS
ONE" MAKES SENSE, BUT AS THEY
SAY, YOU CAN'T ADD GRAPES AND
REINDEER.



EXCEPT IN MY GRAPE AND REIND EER STEW... NUMBERS ARE ALSO UNIQUE IN THAT YOU "DO THEM" ON YOUR FINGERS, WHILE OTHER PARTS OF LANGUAGE HAPPEN MAINLY IN YOUR HEAD... YES, COUNTING HAS BEEN * FROM THE BEGINNING! HOW MANY DAYS IN A MONTH? THREE, FOUR FINE, & WHILE I'M SURE THIS QUESTION HAS AN ANSWER. THE CUPPENT SEEMS INADEQUATÉ TO THE TASK ... THESE GENIUSES CERTAINLY STRAIN THE OLD SOCIAL



AROUND THE SAME TIME, CAVE DWELLERS WERE BEGINNING TO STORE ANOTHER KIND OF INFORMATION AS WELL:
THEY PAINTED REALISTIC ANIMALS ON THE WALLS OF THEIR CAVES — NO ONE KNOWS WHY!



SEVERAL THOUSAND YEARS LATER, THE SUMBRIAINS
DEVISED A SYSTEM TO REPRESENTING THEIR ENTIRE
LANGUAGE IN PICTURES:

CALL IT
"VISI-TALK!"

AND SO WRITING WAS BORN!

UNTIL SOMEONE CAN IMPROVE ON LANGUAGE ITSELF, WRITING WILL BE THE ULTIMATE HUMAN SYSTEM OF INFORMATION STORAGE. It'S HEARLY UNIVERSAL! PEOPLE ALL AROUND THE WORLD INVENTED SYMBOL SYSTEMS TO ENCODE SPOKEN LANGUAGE. OF COURSE, TECHNIQUES VARIED FROM PLACE TO PLACE...

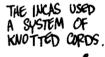


THE
SUMERIANS
WROTE ON
CLAY, TABLOSS,
WHILE THE
EGYPTIANS
USED SOFT
PAPYRUS.



CHIVESE WRITING OBGAN WITH MESSAGES TO THE GODS INKED ON TORTOISE SHEUS.











GREAT!
NOW THAT WE'VE
STORED ALL THAT
INFORMATION, HOW
DO WE FIND
IT AGAIN?

WE'LL RETURN TO THAT POINT LATER!



ALL THE EARLY CIVILIZATIONS HAD WAYS OF REPRESENTING NUMBERS THAT WERE FAR ADVANCED OVER THE STONE AGE TALLY BONE, ON WHICH THE NUMBER IS SIMPLY MADE BY PILING UP 1's. NOT TOO USEFUL...

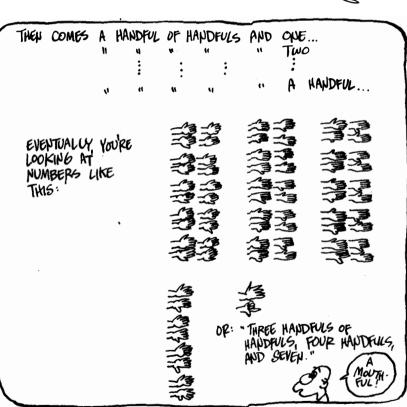


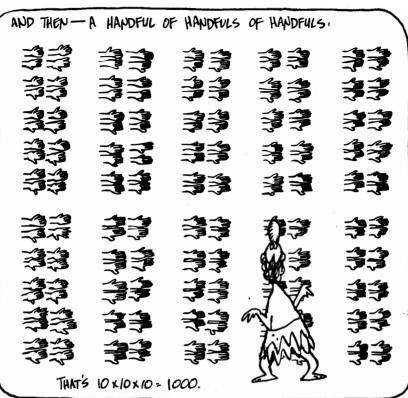
SOMETIME BETWEEN
TALLY BONE AND
CIVILIZATION, REOPLE
BECAME ACCUSTOMED
TO COUNTING BY
FIVES AND TENS—
FOR AN OPPIOUS
REASON: IT WAS
HANDY.

Mm

少辛季

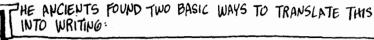
LET'S CALL AFTER A WHILE. TEN A "HANDFUL" YOU REAGI AND DO SOME COUNTING. HANDFUL FIRST COME OF HANDFULS SOME NUMBERS LIKE (TEN TENS, "TWO HANDFULS OR A HUNDRED). AND THREE."





NEXT COMES
TEP THOUSAND...
A HUNDRED THOUSAND...
A THOUSAND THOUSAND...
TEN THOUSAND THOUSAND...
EACH OF WHICH IS
A HANDFUL OF
THE ONE BEFORE!





JUST PILE THEM UP:

TWO HUNDRED THREE TEN SIX ONE NINE
TROUSANDS TEN UNITS

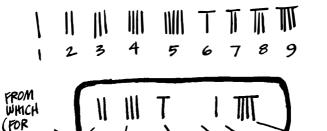
0R 236,019

ASIDE FROM HAVING A CERTAIN GRAPHIC CHARM, THESE NUMERALS ARE VERY EASY TO READ, ONCE YOU'RE USED TO THEM (JUST AS "3 BILLION" READS QUICKER THAN





ON THE OTHER HAND, THE CHINESE USED THE **POSITION** OF NUMERALS TO INDICATE THEIR VALUE. FIRST THEY COUNTED FROM ONE TO MINE:



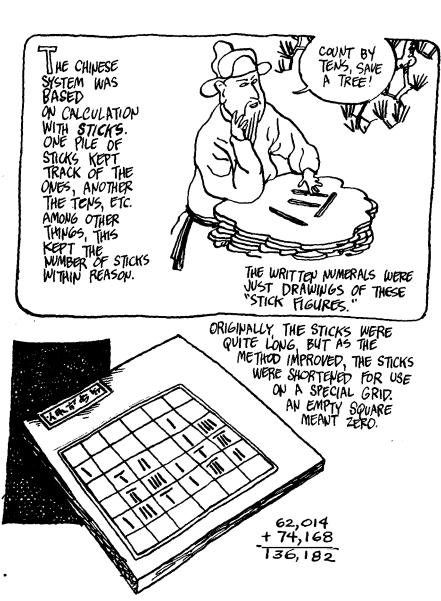
(FOR EXAMPLE):
Two hundred three six no one nine thousands thousands thousands thousands ten units or 236,019.

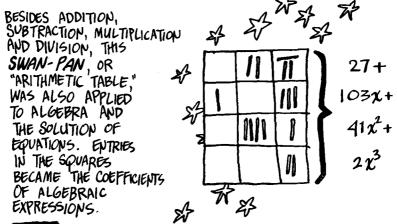
EXCEPT FOR THE UNFAMILIAR NUMERALS, THIS SYSTEM IS NEARLY THE SAME AS OUR OWN.

THE ONLY DIPPERENCE IS THAT IT LACKED A SYMBOL FOR ZERO. THE CHINESE JUST LEFT A BLANK INSTEAD.



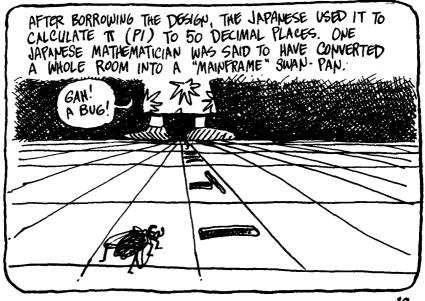
IN PRACTICE, THIS WAS MUCH LESS OF A PROBLEM THAN IT MIGHT HAVE BEEN, BECAUSE THE CHIPESE DID NOT CALCULATE ON PAPER !!!

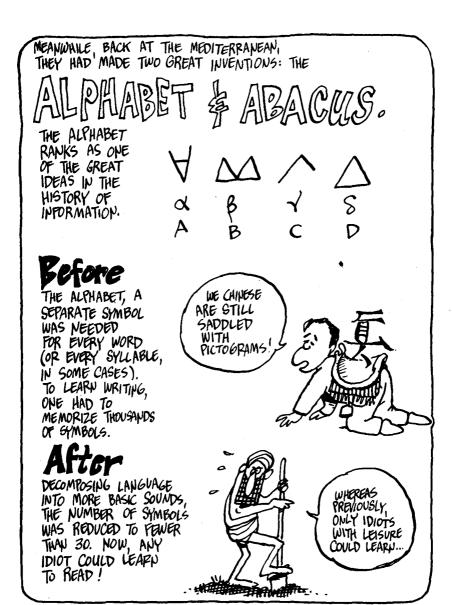






THIS TECHNIQUE HAD THE PICTURESQUE NAME OF "THE METHOD OF THE CELESTIAL ELEMENT."





THERE'S A LESS OBVIOUS ADVANTAGE OF THE ALPHABET, BUT NO LESS IMPORTANT:

alphabetical order.

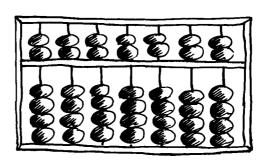


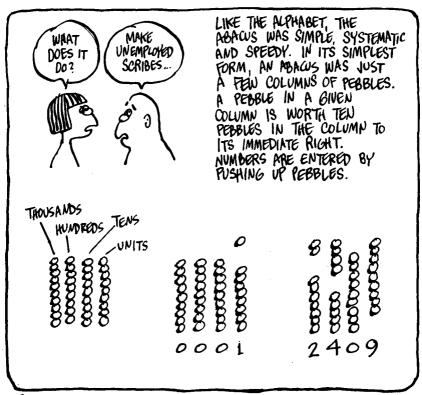
WITH THOUSANDS OF PICTOGRAMS, ANY FILING SYSTEM HAS TO BE COMPLICATED, BUT GIVEN THE ORDER OF AN ALPHABET, YOU CAN PUT WORDS IN ORDER, TOO. IMAGINE USING A PHONE BOOK, DICTIONARY, OR LIBRARY WITHOUT ALPHABETICAL ORDER!



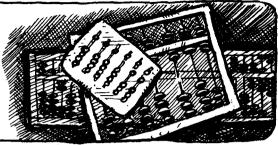
COMPUTERS SPEND A GOOD PAPT OF THEIR TIME JUST PUTTING THINGS IN ORDER!

THE ABACUS, ORIGINALLY 'A PRODUCT OF THE MIDDLE EAST, IS A FULL-BLOWN HAND-HELD DECIMAL CALCULATOR.





THE ABACUS
HAS SEEN MANY
INCARNATIONS
AND BEEN
USED IN MOST
PARTS OF
THE OLD WORLD.

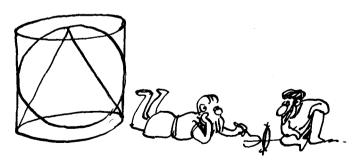


WE KNOW
FROM PICTURES
THAT THE
ANCIENT GREEKS
HAD THE
ABACUS, BUT
THEIR
MATHEMATICIANS
NEVER DISCUSSED
IT. (GREEK
INTELLECTUALS
LOOKED DOWN
ON THE
WORK OF THE
HANDS...)



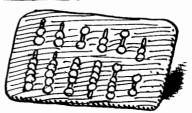


This may have been why greek mathematicians concentrated on geometry...



The Romans

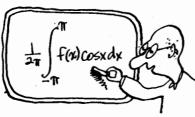
ALSO USED THE ABACUS.
THEIRS CONSISTED OF
MARBLES SLIDING ON A
GROOVED BRONZE PLATE:



IT CONTRIBUTED A COUPLE OF MATHEMATICAL WORDS TO ENGLISH:

IN LATIN,

MEANT LIMESTONE OR MARBLE...



PROM WHICH COMES "CHALK!"

CALCULUS

WAS AN ABACUS PEBBLE... AND DOING ARITHMETIC WAS

CALCULATION

THE
ROMANS DID
NOT CALCULATE
WITH
ROMAN
NUMERALS!!



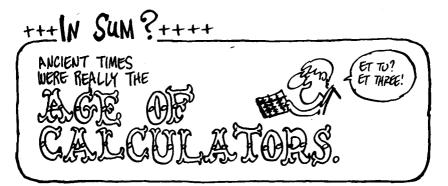
BEATS ME! I LOST MY MARBLES...





...AND FELL...
ROME WAS SACKED...
CHRISTIANITY ROSE
FROM ITS
ASHES... CLASSICAL
LEARNING VANISHED
IN THE WEST...
AND ONLY A FEW
MATH PROBLEMS
REMAINED
LEGITIMATE, LIKE
COMPUTING THE
DATE OF EASTER...
OR HOW MANY
ANGELS FIT ON
THE HEAD OF
A PIN...





ALTHOUGH THE ANCIENTS HAD WAYS OF WRITING NUMBERS, THEY RARRLY CALCULATED IN WRITING.

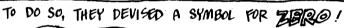
THIS IS NOT SO EASY TO APPRECIATE FOR THOSE OF US WHO WERE RAISED ON PENCIL AND PAPER.

SO THE NEXT HOW CAN WE REMEMBER OUR MULTIPLICATION THOLES?

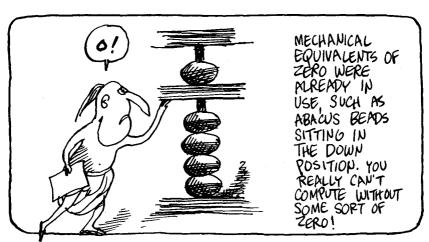
COMPLAIN THAT ELECTRONIC TO CALCULATORS ARE RULVING ARITHMETIC...

... SIMPLY REPLY THAT REOPLE SURVIVED WITH CALCULATORS FOR MORE THAN 4000 YEARS!!

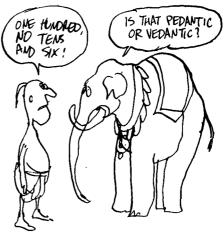








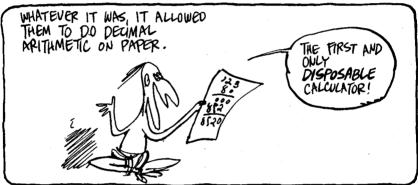
WHY HADN'T ANYONE
THOUGHT OF PUTTING
IT IN WRITING
BEFORE? MAYBE
BECAUSE WRITING
WAS FOR REPRESENTING
SPOKEN LANGUAGE,
AND NOBODY SAYS—



BUT FOR SOME REASON, THE HIPDUS INVENTED A WRITTEN ZERO!

1238486680





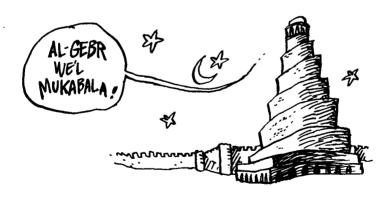
AND SO BEGAN THE AGE OF PENCIL AND PAPER, A MERE 1300 YEARS AGO — PRETTY BRIEF COMPARED WITH THE AGE OF CALCULATORS!!



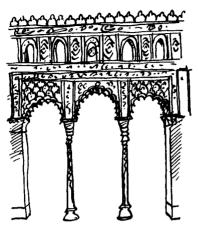
THE INDIAN MATH WAS PICKED UP BY THE AIR AIR S, WHO SPREAD IT AU.
THE WAY WEST TO SPAIN.



ARDUND THE YEAR 830, A PERSIAN SCHOLAR WROTE THE STANDARD TEXT BOOK ON THE SUBJECT. HIS NAME WAS AS AS A PERSIAN SCHOLAR WROTE THE MOHAMMED IBN MUSA ABU DJEFAR, BUT HE WAS KNOWN BOOK?



OF ALGEBRA, FOR SHORT.



BY THE 1100'S, MUSLIM CIVILIZATION HAD GROWN SO MAGNIFICENT THAT THE EUROPEANS WERE BEGINNING TO WONDER...



A FEW INTREPID INFIDERS
WENT TO LIVE AMONG THE ARABS,
LEARNED THEIR LANGUAGE,
SNUCK INTO THEIR UNIVERSITIES,
AND TRANSLATED THEIR
CLASSICS INTO LATIN.



IN AL KHWARISMI'S BOOK THEY FOUND THE INDIAN NUMBRAIS.



AL-KHWARISMI AL-KARISMI ALGARISMI ALGORISMI

PROPOUNCED OFTEN ENOUGH, THE MATHEMATICIANS NAME WAS TURNED INTO

ALGORISM-

WHICH IS WHAT THE EUROPEANS CALLED THE NEW SYSTEM OF CALCULATION.



YETH, VERY NITHE...

FROM THE SAME ROOT COMES

ALGORITHM,

A COMPUTER WORD WE'LL EXPLORE IN A BIT...

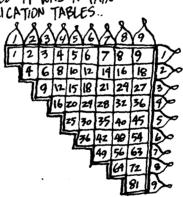


You CAN TURN THIS "O" INTO 6 OR 9... 3 LOOKS TOO MUCH LIKE 8, ETC..

THAT'S WHY I LIKE IT!

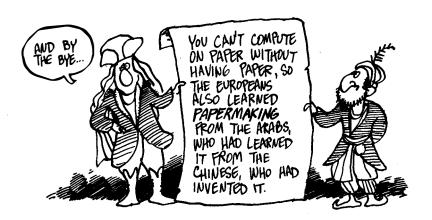
... AND EVERYONE AGREED IT WAS A PAIN TO MEMORIZE MULTIPLICATION TABLES...





BUT IT DID CATCH ON—
NOT NECESSARILY BECAUSE IT
WAS FASTER THAN THE
ABACUS—IT WASN'T—BUT
BECAUSE, AS THE ARABS
KNEW, IT ENCOURAGES
ABSTRACT SYMBOL-MANIPULATION:
FIRST ALGEBRA, AND LATER
THE CALCULUS AND ALL
OTHER HIGHER MATHEMATICS.

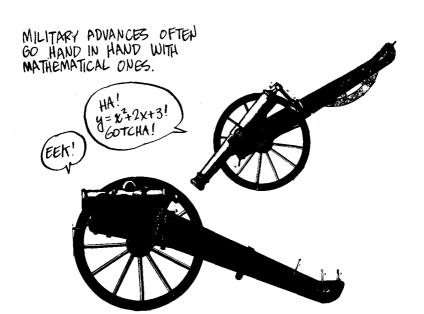




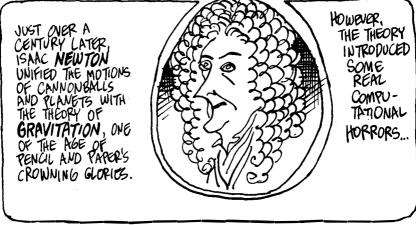
IN EXCHANGE, THE CHINESE TOOK THE ABAC 1 THUK WE CAN MAKE IT WITH ONE AND RAPIDLY MADE IT THEIR LESS BEAD... NO. 1 CALCULATOR. FROM CHINA THE ABACUS SPREAD TO JAPAN, WHERE - MEED I SAY IT? -ITS DESIGN was improved! BUT BACK TO ALGORISM ... 1 WHILE EUROPEAN
SCHOLARS WERE
TRANSLATING THE CLASSICS
IN ARAB LIBRARIES,
THE CRUSADERS
WERE DOING THEIR
BEST TO DESTROY
ISLAMIC CIVILIZATION.





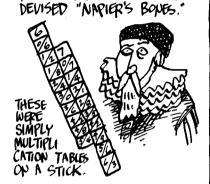


IN THE 1500'S, NICCOLO TARTAGLIA (1491-1559) COMPUTED THE PATHS OF CANNONBALLS (AN IMPORTANT PROBLEM IN THE LATER HISTORY OF COMPUTERS, AS WE'LL SEE).









THE FIRST REAL MACHINE WAS BUILT BY WORKER TO 1592-)

IT COULD ADD, SUBTRACT, MULTIPLY, AND DIVIDE...

MULTIPLY, AND DIVIDE...
BUT WAS LOST IN THE
30-YEARS WAR.
SCHICKARD
HIMSELF

Hingelf Died of Plague And Covlon't

DEPEND HIS PRIORITY,

Blaise PASCALIPE USUALLY GETS CREDIT FOR BUILDING THE FIRST CALCULATOR. HIS "PASCALIPE" COULD OPLY MOD AND SUBTRACT.

6 officed Wilhelm

LETBINTZ (1646-1716)
IMPROVED PASCAL'S DESIGN

DREAMED
OF A
DAY WHEN
ALL
REASONING
COULD BE
DONE BY
TURNING A

QUITE A BIT ...

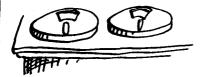
à۷D



DURING THE 1700'S,
MORE SUCH MACHINES
WERE BUILT, BUT
ALL FELL FAR
SHORT OF BEING
ANYTHING LIKE
A GENERAL-PURPOSE
COMPUTER.



FOR EXAMPLE: IN EVERY CASE, THE USER ENTERED NUMBERS BY SETTING A ROW OF WHEELS OR KNOBS ...



... AND THEN
TURNED THE
APPROPRIATE
CRANK TO ADD
OR MULTIPLY.





THE STED ONLY OF THE MUMBERS TO BE COMBINED.

AS WILL BE PLAIN SOOP ENOUGH, AN ALL-PURPOSE COMPUTER MUST ALSO BE ABLE TO DO MORE: IT MUST READ WITH TROSE NUMBERS!

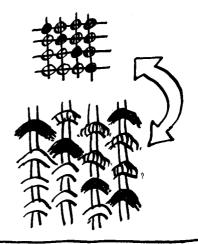


WELL, I
WAS ONLY
TRYING TO
MAKE AN
ADDING
MACHINE...

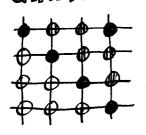
THE GERM OF THIS IDEA CAME NOT FROM THE LAB OR A SCIENTIST'S STUDY, BUT THE SOOTY FACTORIES OF THE



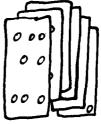
YOU MAY NEVER
HAVE THOUGHT OF
A WEAVER'S LOOM
AS AN INFORMATION
PROCESSOR, AND
YET: IT TRANSLATES
AN ABSTRACT
DESIGN INTO A
PATTERN OF COLORS,
CREATED BY
LOOPING OVER EACH
COLORD THREAD
AT THE
APPROPRIATE PLACE.



IN THE MID-1700'S, A SYSTEM WAS INVENTED FOR REPRESENTING THESE PATTERNS ON PUNCHIED CAIRDS.







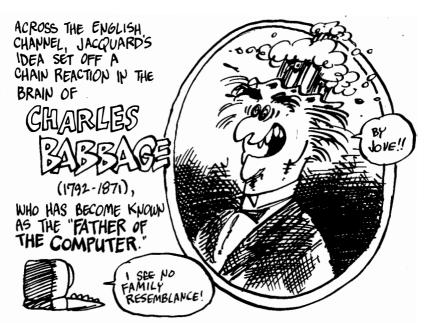
WITH AN OLD-FASHIONED HANDLOOM, THE WEAVER READ THE CARDS, BUT IN 1801, JOSEPH MARIE TALCOMULTH A POWER LOOM WITH AN AUTOMATIC CARD READER.

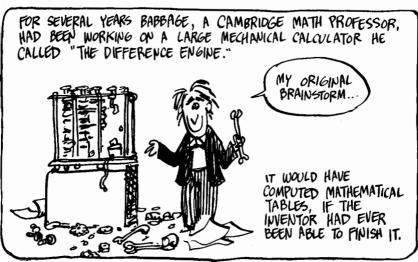




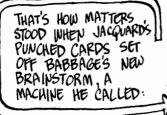
IN WENT THE CARDS, OUT CAME THE CLOTH ...













BECAUSE IT SO CLOSELY RESEMBLED A COMPUTER, LET'S TAKE A CLOSER LOOK AT THE ANALYTICAL ENGINE, AS BABBAGE IMAGINED IT. ITS COMPONENTS INCLUDED—

THE MULLS

AT THE ENGINE'S
HEART WOULD BE A
GREAT NUMBERCRUNCHER, AN ADDING
MACHINE ACCURATE
TO 50 DECIMAL
PLACES. BABBAGE
CALLED THIS THE MILL.

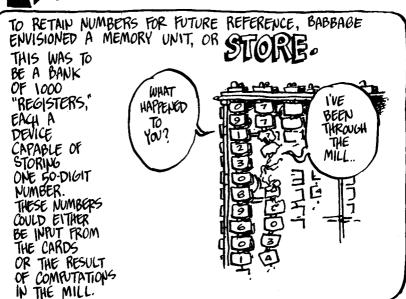




THAT IS, THE PUNCHCARDS CONVEYED NOT ONLY THE NUMBERS TO BE CRUNCHED BUT ALSO THE PATTERN OF CRUNCHING!



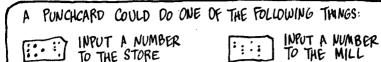
SO THE MACHINE WOULD NEED A SPECIAL CARD-READING INPUT DEVICE.

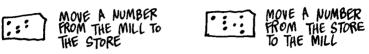




FINALLY, THE

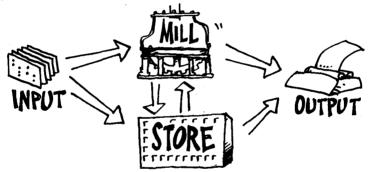
BABBAGE
DESIGNED THE
WORLD'S FIRST
AUTOMATED
TYPE SETTER
TO PRINT THE
RESULTS OF
COMPUTATIONS.

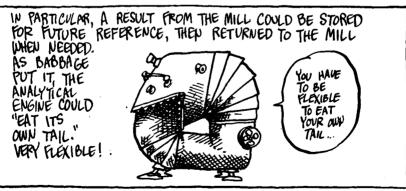




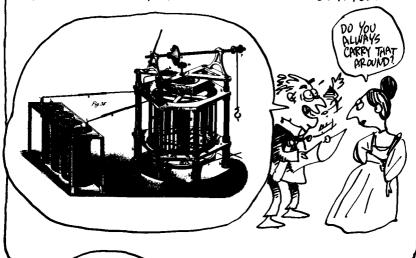


WHICH MAY BE SUMMARIZED IN THIS DIAGRAM:





SO FAR, THESE IDEAS WERE STILL ON THE DRAWING BOARD. NOW BABBAGE BEGAN LOOKING FOR SYMPATHETIC SOULS WHO COULD HELP PUT HIS PLANS INTO OPERATION.





THE MOST SYMPATHETIC WAS

LADY LOVELACE, DAUGHTER OF THE POET LORD BYRON AND AN ENTHUSIASTIC AMATEUR MATHEMATICIAN.
IF CHARLES BABBAGE IS THE COMPUTER'S FATHER, ADA LOVELACE IS ITS MOTHER!

ADA BECAME THE FIRST PROGRAMMER: SHE WROTE OUT ACTUAL SEQUENCES OF INSTRUCTIONS FOR THE ANALYTICAL ENGINE ...



SHE INVENTED THE SUBROUTINGS A SEQUENCE OF INSTRUCTIONS WHICH CAN BE USED AGAIN AND AGAIN IN MANY CONTEXTS. we can



OP THEM

HAVE A WHOLE LIBRAPY

SHE RECOGNIZED THE VALUE OF [GOOP] NOS THERE SHOULD BE AN INSTRUCTION THAT BACKS UP THE CARD READER TO A SPECIFIED CARD, SO THAT THE SEQUENCE IT INITIATES CAN BE EXECUTED REPEATEDLY



AND SHE DREAMED UP THE GONDOTTIONAL JUMAP: THE CARD READER COULD "JUMP" TO ANOTHER CARD 118 SOME CONDITION IS SATISFIED.



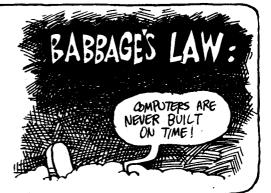
NOT BAD FOR A MACHINE THAT NEVER EXISTED...
THE GOVERNMENT
REFUSED TO SUPPORT IT,
IN VIEW OF BABBAGE'S
TRACK RECORD WITH
THE DIFFERENCE ENGINE.
THEY CALLED IT:

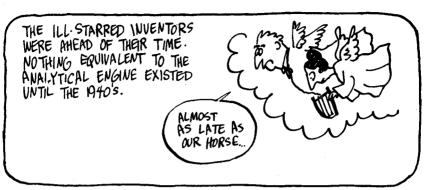


DESPERATE FOR FUNDS, BABBAGE COOKED UP A "SCIENTIFIC" RACETRACK BETTING SCHEME — AND SQUANDERED ADA'S FORTUNE.



THE STORY ENDED
UNHAPPILY: ADA
DIED YOUNG...
AND BABBAGE
NEVER FINISHED
THE ANALYICAL
ENGINE, WHICH
BECAME THE
FIRST EXAMPLE OF—





IN THE MEANTIME, MATTERS PROGRESSED IN TWO DIRECTIONS.

ON THE
ONE HAND
WERE MECHANICAL
CALCULATORS:
SEVERAL
ENGINEERS BUILT
BABBAGEINSPIRED
DIFFERENCE ENGINES.
FOR SOME
REASON, THESE
NEVER CAVENT ON...







...ALTHOUGH DESKTOP ADDING MACHINES AND CASH REGISTERS DID BECOME FIXTURES IN BUSINESS. ON THE OTHER HAND WERE
THE PUNCHCARD MACHINES,
BEGINNING WITH THE CENSUS
TABULATORS DESIGNED BY

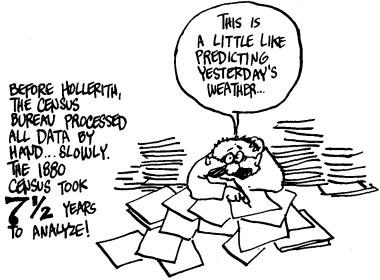
HERMAN

HOLLERITA (1867).

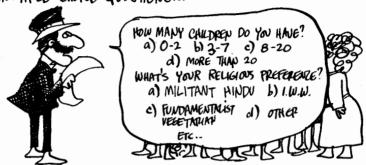
INSPIRED, AS BABBAGE HAD BEEN, BY THE JACQUARD LOOM, HOLLORITH INVENTED A MACHINE PURELY FOR ACCUMULATING AND CHASSIPYING INFORMATION.



BECAUSE THIS WAS A NEW SORT OF JOB FOR A MACHINE— AND THE KIND FOR WHICH COMPUTERS ARE IDEALLY SUITED— LET'S TOOK A CLOSER LOOK.



THEN AS NOW, THE CENSUS FORM CONSISTED OF A SERIES OF MULTIPLE CHOICE QUESTIONS...



FROM THIS, ONE WANTED TO FIND: G

THE TOTAL NUMBER OF CITIZENS...

HOW MANY HAD 0-2 CHILDREN...

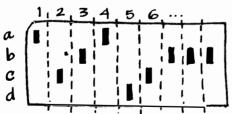
HOW MAPY WERE MILITANT HINDUS...

AS WELL AS SUCH THINGS AS:



HOW MANY
PULIDAMENTALIST
VICETAPIANS
HAVE MORE
THAN 20
CHILDREN 9

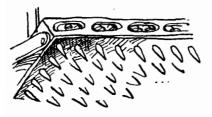
HOLLERITH PROPOSED TO PUT EACH PERSON'S RESPONSES ON A SINGLE PUNCHED CARD THE SIZE OF AN 1880 DOLLAR BILL. TO OVER-SIMPLIFY SLIGHTLY



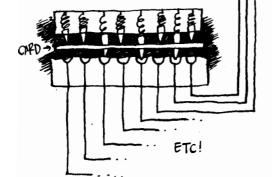
SIMPLIFY SLIGHTLY, EACH COLUMN, PEPRESENTED ONE QUESTION. THE HOLE IN A GIVEN COLUMN INDICATED THE ANSWER TO THAT QUESTION.

THIS CARD SHOWS RESPONSES OF 1-a, 2·c, 3·b, 4-a, 5-d, ETC...

THE CARDS WERE "READ"
BY A DEVICE CONSISTING
OF A GRID OF LITTLE
PINS MOUNTED ON
SPRINGS AND WIRED
ELECTRICALLY.

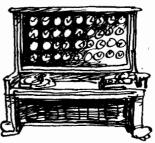


WHEN BROUGHT INTO CONTACT WITH THE CARD, ONLY THOSE PINS LYING OVER A HOLE WOULD PASS THROUGH. EACH OF THESE DIPPED INTO A SMALL CUP OF MERCURY, COMPLETING AN ELECTRICAL CIRCUIT.





EACH CUP WAS WIRED TO A COUNTER, WHICH ADVANCED EACH TIME AN ELECTRIC PULSE ARRIVED.



AND SO THE RUNNING TOTALS OF EVERY POSSIBLE RESPONSE WERE CONTINUOUSLY DISPLAYED!



THE TABULATOR ALSO HELPED ANSWER QUESTIONS SUCH AS:
"HOW MANY PEOPLE WHO ANSWERED 2-a ALSO ANSWERED 3-c?"

MEAHING: HOW MAHY MILITANT HINDUS LIVE IN KANSAS?

HERE'S HOW:

FIRST, AFRANGE A BELL TO RING WHENEVER A CAPD WITH 2-A IS ENTERED.



THEN RUN
THROUGH
ALL CARDS,
PULLING OUT
ALL THOSE
THAT RING
THE
BELL.



THIS CREATES A STACK OF ALL THE MILITANT HIPDU CARDS. RUN THESE THROUGH THE TABULATOR A GAMP.



The MACHINE THEN SHOWS ALL THE TOTALS FOR MILITANT HINDUS.



THIS SORT OF
JOB — AHALYZING
AND COMPARING
LARGE AMOUNTS
OF WEORMATION—
IS NOW KNOWN
AS:



THE HOLLERITH TABULATOR CUT THE DATA PROCESSING
TIME FOR THE 1890 CENSUS
BY THOM THINKDS, TO
212 YEARS. THIS SOUNDS
LONG NOW, BUT AT THE
TIME, IT WAS IMPRESSIVE!!



HOLLERITH FOUNDED A COMPANY TO MANUFACTURE HIS CARD-OPERATED DATA PROCESSORS, AND HE FOUND A NUMBER OF TAKERS:

A RAILROAD COMPANY USED THE TABULATOR FOR AUDITING PREIGHT STATISTICS ... A TOOL MANUFACTURER TURNED IT TO COMPILING COSTS. ANALYZING PAYROLL, AND MANAGING INVENTORY ... A WHOLESPLE HOUSE NEEDED IT TO KEEP TRACK OF MERCHAMDISE, SALES SALESMEN, CUSTOMERS, ETC ETC ETC ...



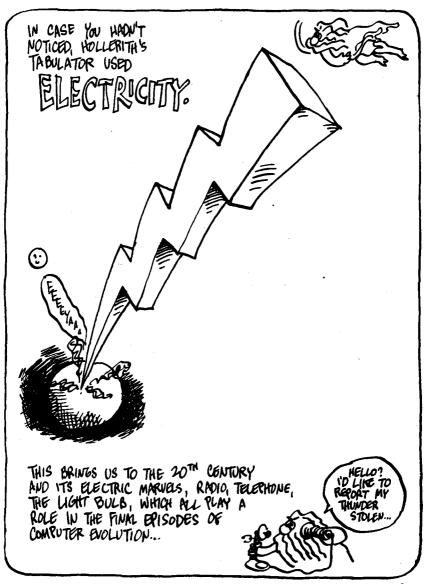
THIS COMPANY

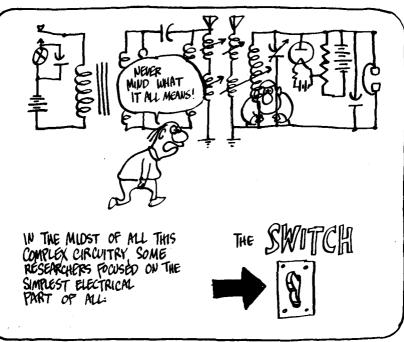
HOLLERITH'S COMPANY DID FAIRLY WELL ... LATER, IT
GOT INTO COMPUTERS, TOO... AND DID WELL ... YOU
MAY HAVE HEARD OF IT... TODAY IT'S CALLED

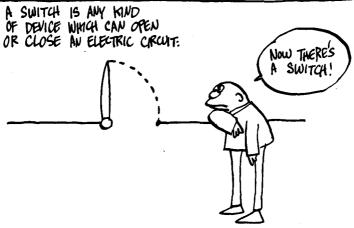
SO BIG, IT DOESN'T FIT IN THE PANEL!

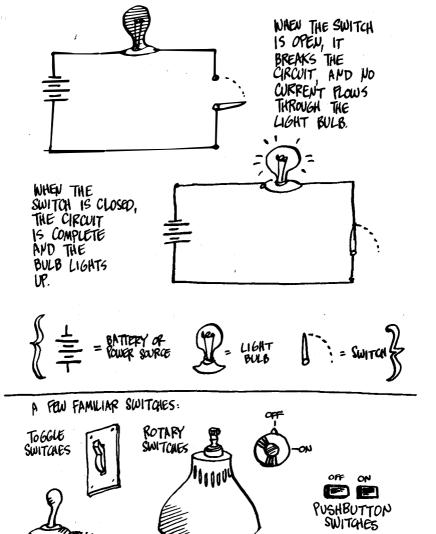


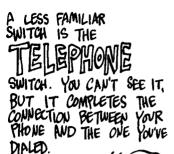










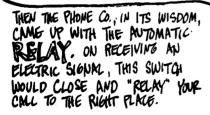


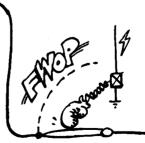


YES ...
UH HUH...
SO TRUE...
AM HAM
YES YES...

IN THE OLD DAYS, THIS HAD TO BE DONE BY HAND-

THE OPERATOR'S WORK STATION WAS CALLED A SIMPLICATIONARD, AFTER NU!





THE TELEPHONE RELAY COULD SWITCH MUCH FASTER THAN THE HUMAN HAND — ABOUT 5 TIMES PER SECOND! IT MADE THE SWITCHBOARD OPERATOR OBSOLETE...



GUESS I'LL APPLY FOR A JOB AT THE RELAY WORKS...



BUT IT COULON'T HOLD A CANDLE TO ANOTHER TYPE OF SWITCH INVENTED BEN EARLIER:
THE VACUUM
TO UBB.



REMEMBER WHEN
TUBES
USED TO
GLOW IN THE
BACK OF THE
RADIO?
YOU DON'T?
SIGH....



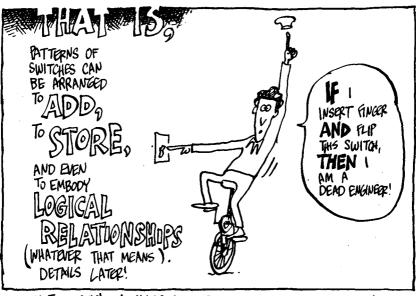
THE TUBE CAN ALSO BE FLIPPED ON AND OFF LIKE A SWITCH, SO FAST YOU CAN'T EVEN SEE IT FLICKER: IT JUST GLOWS... BUT IT CAN SWITCH AS OFTEN AS

1,000,000 TIMES PER SECOND!!!





NOT LONG AFTER THESE SWITCHES WERE INVENTED, PEOPLE REALIZED THAT THEY COULD BE COMBINED INTO COMPUTER COMPONENTS!







Who built

THE FIRST ELECTROMECHAMICAL COMPUTER? THE VERY FIRST WAS TECHNOLOGY

7033 (1910 -).

MS Z.1, BUILT IN 1936, CALCULATED WITH RELAYS AND READ INPUT FROM PUNCHED FILM.



2USE, A GERMAN, TRIED TO SELL THE 2.1 TO HIS GOVERNMENT FOR WAR WORK.



THE NAZIS ASSUMED THEY HAD "ALL BUT" WON THE WAR, SO THEY TURNED HIM DOWN... AND POSSIBLY CHANGED HISTORY!!

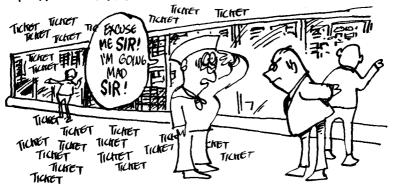




IN THE USA, THE NAVY COLLABORATED WITH HARVARD AND IBM TO CONSTRUCT THE MARK B, AN ELECTROMAGNETIC GIANT LAUNCHED IN 1944



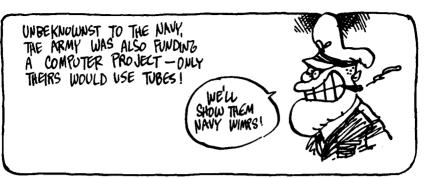
DESIGNED BY HARVARD PROF HOWARD AIKEN, WHO MODELED IT ON BABBAGE'S ANALYTICAL ENGINE, MARK I OCCUPIED SOME 1200 CUBIC FEET AND CONTAINED THOUSANDS OF RELAYS. WHEN IT CRANKED UP, THEY SAY IT SOUNDED LIKE A MILLION KNITTING NEEDLES!



MARK I COULD MULTIPLY TWO 10-DIGIT NUMBERS (A CONVENIENT MEASURE OF COMPUTER SPEED) IN ABOUT

3 seconds.





THEIR AIM WAS THE SAME AS TARTAGUA'S IN THE 1500'S: TO COMPUTE BALLISTICS MORE ACCURATELY.

TAPTAGUA HAD ERRED IN SAYING THAT CANNON-BALLS FLY IN PARABOLIC PATHS. IN REALITY. AIR RESISTANCE ALTERS THEIR TRAJECTORY APPRECIABLY. A W GYA VERY COMPLEX WAY, BECAUSE AIR RESISTANCE DIMINISHES AT MOHER AUTITUDES.



IN WORLD WAR I, THE BERMAN
CANNON "BIG BERTHA" SHOT 94 MILES—
TWICE AS FAR AS EXPECTED PROM
OVERSIMPLIFIED CALCULATIONS!

GUNNERS AND
BOMBARDIERS
THEREFORE NEEDED
ACCUPATE
BALLUSTIC
TABLES TO
AIM BY. THESE
COULD HARDLY BE
CALCULATED
ON THE PLY!



BALLISTIC TABLES USED TO BE CALCULATED BY ROOMFULS OF "GIRLS" WITH ADDING MACHINES — AND EVEN THIS WAS SLOW.





THE RESULT OF THEIR LABORS WAS THE BARN SIZED ENLACED INTEGRATOR AND CALCULATOR. WITH 18,000 TUBES, ENIAC WAS FAST:







SO THE ARMY PUT ENIAC TO WORK OP THE NEXT WAR, DOING CALCULATIONS FOR THE NUCLEAR WEAPONS PROGRAM...



BUT STILL IMPRESSIVE:
WITH 18,000 TUBES
FLICKERING ON
AND OFF 100,000
TIMES PER SECOND,
ENIAC HAD TO
PERFORM FAR MORE
RELIABLY THAN
NY MACHINE ENER
CONSTRUCTED.

MISCALCUATION
COULD BE
FATAL!

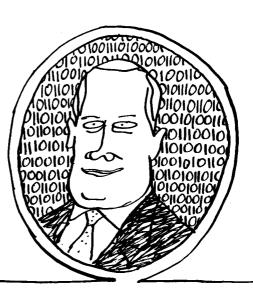
NOW ENTERS

TOMN VOR

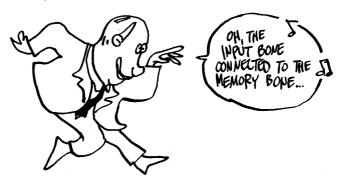
NEUMANN
(1903-1957), A
PRINCETON MATH
PROPESSOR WHO
MORE THAN ANYONE
GETS CREDIT FOR
TURNING ELECTRONIC
CALCULATORS

INTO "ELECTRONIC

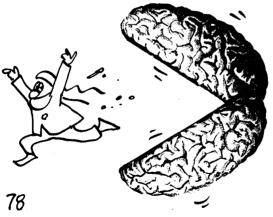
BRAINS.



VON NEUMANN PONDERED THE COMPUTER'S LOGICAL STRUCTURE IN THE MOSTRACT: HOW IT CONTROLS ITSELF, HOW MUCH MEMORY IT NEEDS AND WHAT FOR, ETC... AND HE ASKED HIMSELF HOW COMPUTERS COULD BE MADE MORE LIKE HUMAN "WIRING," I.E., THE CENTRAL NERVOUS SYSTEM.







YOUR BRAIN IS FULL OF THESE "STORED PROGRAMS":
YOU KNOW HOW TO TIE YOUR SHOELACES, HOW TO FEED YOURSELF, HOW TO MULTIPLY
94 TIMES 16, HOW TO TALK, HOW TO WALK...

VON NEUMANN PROPOSED TO MAKE COMPUTERS DO LIKEWISE:

II.

FIND A WAY TO ENCODE THE INSTRUCTIONS
INTO A FORM WHICH COULD BE STORED IN
THE COMPUTER'S MEMORY. VON NEUMANN SUGGESTED
USING STRINGS OF ONES AND ZEROS.







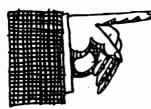
Do

STORE THE INSTRUCTIONS IN MEMORY, ALONG WITH WHATEVER OTHER INFORMATION (NUMBERS, ETC) IS NEEDED TO DO THE PARTICULAR JOB.





WHEN RUNNING THE PROGRAM, FETCH THE INSTRUCTIONS STRAIGHT FROM MEMORY, RATHER THAN READING A NEW PUNCHCARD AT EACH STEP.



THIS IS THE CONCEPT OF THE STORED PROGRAM

The advantages?



SPEED O LIKE THE SURGEON,
THE COMPUTER FINDS IT
MUCH FASTER TO WHIZ
INSTRUCTIONS FROM "BRAIN"
TO "FINGERS" THAN TO
"RETURN TO THE TEXTBOOK"
AFTER EXECUTING EACH STEP.



WITH SEVERAL PROGRAMS STORED AT ONCE, THEY CAN REFER TO ONE ANOTHER RUMMING IN COMBINATION. SURGERY IS ACTUALLY SUCH A COMBINATION.



SELF MODIFICATION IS STORED ELECTRONICALLY, PROGRAMS MAY EASILY BE WRITTEN WHICH CAN MODIFY OR ADJUST THEMSELVES.
THIS TURNS OUT TO BE CRITICALLY IMPORTANT!



TO MAKE HIS POINT, VON NEUMANN WROTE SOME CODE FOR A PROGRAM CALLED:





IT'S A SIMPLE JOB TO DESCRIBE:

GIVEN TWO LISTS OF NAMES (FOR EXAMPLE):

ALABAMA, S. ANTEATER, J. ANTEATER, B. AARDVARK, A. TARDIGRADE, C. BEAVER, M. OWL, H. ALUGATOR, A.



AARDVARK, A.
ALABAMA, S.
ALLIGATOR, A.
ANTEATER, B.
ANTEATER, J.
BEAVER, M.
OWL, H.
TARDIGENDE, C.

MAKE ONE LIST IN ALPHABETICAL ORDER.

THIS SECTIONALLY SIMPLE PROCESS BECOMES HORRIBLY TIME-CONSUMING WHEN THE LISTS ARE LONG.

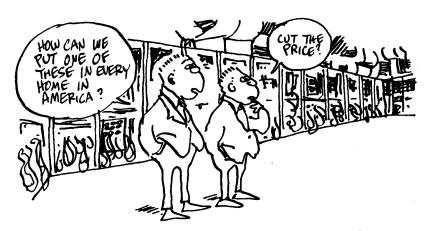
SOO:
HERE'S ANOTHER
IDEAL COMPUTER JOB
THAT CONTAINS
ESSENTIALLY NO
MATH. YOU CAN
SEE HOW THIS ONE
MIGHT APPEAL TO
SOMEONE COMPILING
A TELEPHONE
DIRECTORY OR A
MAILING LIST!!





www.www.www.

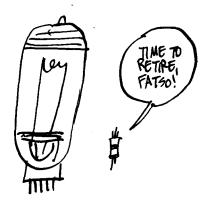




IF COMPUTERS HAD REMAINED AS BULKY AS **ENIAC**, THEY WOULDN'T BE WHAT THEY ARE TODAY... BUT THEY DIDN'T, AND THEY ARE...

IN 1947, THE YEAR
AFTER ENIAC WAS
PINISHED, A TEAM AT
STANFORD INVENTED THE
TIRANSISTERS
USING ELEMENTS CALLED
SEMICONDUCTORS.
LIKE TUBES, TRANSISTORS
CAN ACT AS SWITCHES,
BUT THEY RE

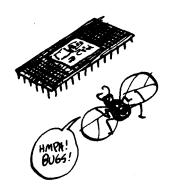
SMALLER,
PASTER,
COOLER, AND
LONGELLINED,
AND THEY DRAW
FAIL LESS
ELECTRIC POWER.



THE FIRST TRANSISTOR12ED COMPUTERS WERE AND SO ROOM-SIZED, NOT BARN-SIZED, AND THER "COMPUTER EFFOR ENTERED EVERYDAY COST (A COUPLE OF LIFE MILLION DOLLARS) WAS AFFORDABLE BY LARGE BUSINESSES AND 10056 100 UNIVERSITIES.

THEN THE TRANSISTOR BEGAN TO SHOW AN INCREDIBLE ABILITY TO SHRINK IN SIZE AND PRICE.

PIRST CAME A WHOLE BOARDPUL OF TRANSISTORS MANUFACTURED AS A SINGLE UNIT ... THEN SI AND VSLI), WHICH DALKED HUNDREDS OF THOUSANDS OF TRANSISTORS ON A





TINY CHP!

AS COMPONENTS GHRANK, THE INDUSTRY EXPLODED!

ANNICOMPUTER
APPEARED. IT WAS THE
SIZE OF A DESK!



IN THE 'TO'S CAME THE MICRO, WHICH CAN BE AS SMALL AS YOU LIKE.



BY THIS TIME, BIG COMPUTERS, ALSO KNOWN AS

MAINFRAMES, HAD BECOME IMMENSELY POWERFUL.



SUPERCOMPUTERS,
WHICH CALCULATE AT RATES

WHICH CALCULATE AT RATES UP TO 500 MEGAPLOPS* — A MILLION TIMES FASTER TRAN ENIAC!

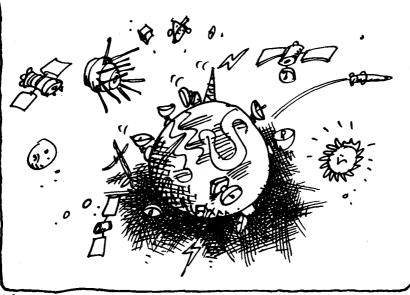


* MILLION FLOATING POINT OFFICERS

THERE'S NO END IN SIGHT.... NOW WE HAVE MICROS WITH THE POWER OF MINIS, "SUPERMINIS" THAT RIVAL MAINFRAMES, MINIS ON A CHIP... AND THERE'S TRUK OF REDUGNO COMPONENTS TO MOLECULAR SIZE USING RECOMBINANT DNA TECHNOLOGY...



THERE SEEMS TO BE NO SUCH THING AS A COMPUTER WITH TOO MUCH COMPUTING POWER. NO MATTER THE SPEED OR CAPACITY, COMPUTERS ALWAYS PIND JOBS TO DO... AND NO WONDER: THIS IS THE AGE OF EXCESS INFORMATION!



PART I

LOGICAL SPAGHETTI





HOW DOES ONE GET TO THE HEART OF THE MATTER?





IF THERE'S ONE
IDEA WE'VE TRIBO
TO, DRUM IN,
IT'S THAT THE
COMPUTER IS
ESSENTIALLY AN
INFORMATION
PROCESSOR.
SO, FORGET THE

TO UNDERSTAND INFORMATION PROCESSING, IT HELPS TO COMPARE IT WITH A MORE FAMILIAR PROCESS: COOKING. SO STEP INTO CRANDMOTHER BABBAGE'S KITCHEN, AS SHE PREPARES BASIC SPACHETTI...

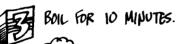


HERE'S THE WORLD FAMOUS RECIPE:



















THIS SPAGHETTI IS BETTER ANALYZED TRAY EATEN! IT'S NOT HARD TO DISTINGUISH A FEW COMPONENTS IN THIS PROCESS:

FIRST, THE INGREDIENTS,







NEXT, THE EQUIPMENT WHICH DOES THE COOKING: HANDS, KETTLE, STOVE, SALTSHAKER, SIEVE, PLATE. SPOON.





THESE FORM THE

PROCESSING UNIT.

LESS OBVIOUSLY, THERE IS A PART OF THE COOK'S BRAIN WHICH CONTROLS THE PROCESS. IT MONITORS AND DIRECTS THE STEP-BY-STEP UNFOLDING OF THE RECIPE. THIS IS REFERRED TO AS THE CONTROL UNIT.



AMD OF COURSE THE COMPLETED DISH, OR

OUTPUT.



WHICH ALSO RESEMBLES THE COOK BRMN...

OF COURSE, SPAGHETTI IS
NOTHING SPECIAL! ANY
RECIPE COULD BE PROCESSED
BY THE SAME BASIC
STRUCTURE:









DOUTPUT



OR, MORE ABSTRACTLY:

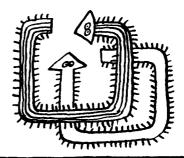


INPUT ⇒ PROCESSING ⇒ OUTPUT

WHITE ARROWS (⇒) ARE THE FLOW OF FOOD
6RAY ARROW (⇒) IS THE FLOW OF INFORMATION
BLACK ARROW (→) IS THE FLOW OF CONTROL.

WITH COMPUTERS, THE DIAGRAM IS SLIGHTLY DIFFERENT:

THERE ARE TWO REASONS
FOR THIS: ONE IS THE
FACT THAT INPUT AND
OUTPUT ARE INFORMATION,
NOT FOOD — SO THE
GRAY ARROW IS THE
SAME AS THE WHITE
ONES.



THE OTHER IS THE GREAT IMPORTANCE OF MISMORY, WHICH FORMS THE FIFTH AND FINAL COMPONENT. IN COMPUTERS, ALL INFORMATION PASSES INTO MEMORY FIRST! HERE'S THE DIAGRAM:

CONTROL PROCESSING INPUT AMEMORY OUTPUT

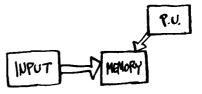
=> = INFORMATION FLOW



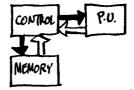


IN THE CASE OF COMPUTERS, THE INPUT CONSISTS OF ALL THE "RAW" DATA TO BE PROCESSED—AS WELL AS THE ENTIRE "RECIPE," OR PROGRAM, WHICH SPECIFIES WHAT'S TO BE DONE WITH THEM.

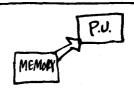
THE MEMORY STORES
THE INPUT AND
RESULTS FROM THE
PROCESSING UNIT:



CONTROL READS
THE PROGRAM AND TRANSLATES
IT INTO A SEQUENCE OF
MACHINE OPERATIONS.



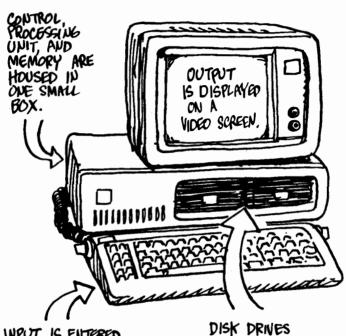
THE PROCESSING UNIT
PERFORMS THE ACTUAL ADDITIONS,
MULTIPLICATION, COUNTING,
COMPARISON, ETC., ON INFORMATION
RECEIVED FROM MEMORY.



THE **OUTPUT** CONSISTS OF THE PROCESSING UNIT'S RESULTS, STORED IN MEMORY AND TRANSMITTED TO AN OUTPUT DEVICE.



HERE'S THE REAL THING (AN IBM PERSONAL COMPUTER), JUST. TO GIVE ONE EXAMPLE OF HOW THESE COMPONENTS MAY ACTUALLY LOOK:



INPUT IS ENTERED PROM KOYBOARD. DISK DRIVES PROVIDE EXTRA MEMORY STORAGE

OTHER COMMON INPUT/OUTPUT DEVICES (NOT PICTURED) ARE A **MODEM**, FOR SENDING AND RECEIVING SIGNALS OVER THE PHONE, AND A **PRINTER**, FOR PRODUCING OUTPUT ON PAPER.

LET'S START IN THE MIDDLE, WITH THE

PROCESSING IIII

IN THE KITCHEN, A CHEF MAY DISPLAY A RICH REPERTOIRE OF PROCESSING POSSIBILITIES:



BUT, AS THE GREAT ESCOFFIER HIMSELF HAS REMARKED, ALL COOKING TECHNIQUES ARE COMBINATIONS OF SIMPLER STEPS: THE APPLICATION OF MORE OR LESS HEAT, WET OR DRY, ETC...





LIKEWISE, ALL THE POWER OF THE COMPUTER DEPENDS ON A COUPLE OF ELEMENTARY OPERATIONS.



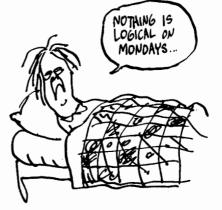
O.K... O.K... NO MORE BEATING AROUND THE BUSH WITH CULINARY METAPHORS...

THE COMPUTER'S ELEMENTARY OPERATIONS ARE





What's A Logical operation, you ask? A Logical question, considering how much easier it is to think of Illogical operations, like amputation of the Thumbs or getting out of Bed on Mondays...



TO EVERYONE'S GOOD FORTUNE, LOGIC ISN'T AS HARD AS IT USED TO BE. IN ARISTOTLE'S TIME, THE SUBJECT WAS DIVIDED INTO INDUCTIVE AND DEDUCTIVE BRANCHES, INDUCTIVE LOGIC BEING THE ART OF INFERRING TRUTHS BY OBSERVING NATURE, WHILE DEDUCTIVE LOGIC DEDUCED TRUTHS FROM OTHER TRUTHS:

- 1. YOU ARE A MAN.
- 2. ALL MEN ARE MORTAL.
- 3. THEREFORE, YOU ARE
 MORTAL.

AHEM :
HOW DO YOU
HOW ALL
MEN ARE
MORTAL ??

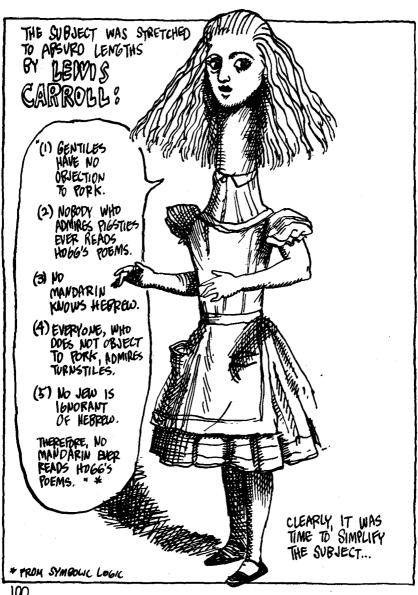
Medieval

LOGICIANS COMPOUNDED
THE COUPUSION
WITH SIX "MODES":
A STATEMENT WAS
EITHER TRUE,
FALSE, NECESSARY,
CONTINGENT,
FOSSIBLE, OR
IMPOSSIBLE.



necessary is to contingent as true is to palse... possibly...

THEIR REASONING GREW
SO MINDLESS THAT
THE MEDICIAL
LOGICIAN DUNS
SCOTUS HAS BEEN
IMMORTALIZED IN
THE WORD
"DUNCE"!



This step was taken by GEORGE

BOOLE (1845. 1864),

AN ENGUSH MATPEMATICIAN WHO BUILT AN "ALGEBRA" OUT OF LOGIC



THAT IS, HE MADE
LOGIC FULLY

SYMBOLIC, JUST
LIKE MATH. SENTENCES
WERE DENOTED BY LETTERS
AND CONNECTED BY
ALGEBRAIC SYMBOLS — AN
IDEA GOING BROWN TO
LEIBNIZ, WHO HAD DREAMED
OF "JUSTICE BY ALGEBRA."

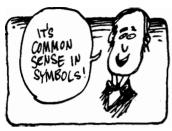


(1-x).(1-y) = 1-x-y+xy. Thefore, 30 Yenps!

WE CAN'T POSSIBLY DESCRIBE BOOLE'S ALGEBRA IN ITS ENTIRETY. WE'LL LIMIT OURSELVES TO THREE WORDS:



BOOLE LOOKED AT THE VERY CONNECTIVE TISSUE OF LANGUAGE: THE WORDS "AMD", "OR", AMD "NOT".





SUPPOSE P IS ANY STATEMENT ... FOR EXAMPLE,

P = "The pig has spots."

ACCORDING TO BOOLE, THIS SENTENCE IS ELTHER TRUE (T) OR FALSE (F). NO OTHER OPTION IS ALLOWED!*



NOW LET Q BE ANOTHER STATEMENT-LIMEUISE TRUE OR FAISE:

Q = "The pig is glad."



NOW FORM THE COMPOUND SENTENCES:

P AND Q = THE PIG IS SPOTTED AND THE PIG IS GLAD.

P OR Q = THE RIG IS SPOTTED OR THE PIG IS GLAD.

WHEN ARE THESE SENTENTES TRUE!

POPQ? F RE

* IN SOME VERSIONS OF LOGIC, MORE THAN TWO TRUTH VALUES ARE PERMISSIBLE.

THERE ARE FOUR POSSIBLE COMBINATIONS OF TRUTH AND FALSEHOOD FOR P AND Q.







P FALSE Q TRUE





P TRUE, Q FALSE P FALSE, Q FALSE



"The PIG IS GLAD AND HAS SPOTS."

THIS IS TRUE ONLY IN THE ONE CASE in which P.O ARE BOTH TRUE. THIS IS SUMMARIZED IN A TRUTH TABLE:



8	Q	P	AND Q
T	T	11	T
T	F	1	F
F	T	11	E
<u> </u>	۴	11	F



"THE PIG IS GLAD OR HAS SPOTS."

THIS IS TRUE IN THE THREE CASES FOR WHICH EITHER ONE OF THE Statements P, Q is true.



P	Q	PORQ
T	T	T
F	Ť	1 +
F	F	F

EXCEPT FOR THE ONE WARD EQUATION

|\Omega| = |, These Look

LIKE ORDINARY

ARITHMETIC ... WITH

"ANO" PLAYING THE ROLE

OF "TIMES" AND "OR"

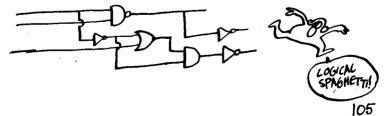
IN THE ROLE OF "PLUS."

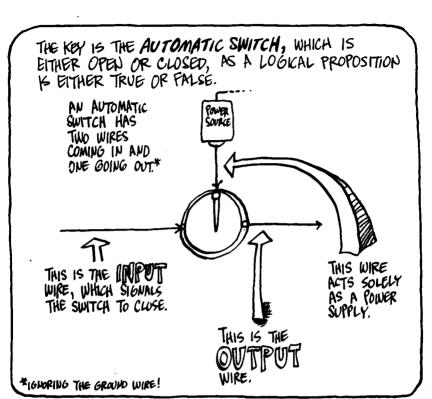


WE'RE NEVER GOING TO USE THE SYMBOLS • AND @... YOU CAN FORCET ABOUT THEM... BUT USING I AND O TO REPRESENT TRUE AND FALSE IS VERY USEPUL... SO FROM NOW ON WE'LL WRITE TRUTH TABLES LIKE THIS:

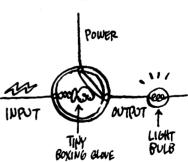
	P 911	PAMOQ	PQI	PORQ	P II NOT-	P_
.	11	1	11	1	1 0	
D	10	0	10	ı	0 1	
8	01	0	01	1	•	
	0011	0	00	0		

FROM THESE RELATIONATIONS, BOOLE DUILT UP AN ENTIRE ALGEBRA, USING ONLY THE NUMBERS O AND 1... TOONY THIS BOOLEAN ALGEBRA IS USED ALL THE TIME BY COMPUTER ENGINEERS — ONLY THEY EXPRESS IT AS ELECTRICAL CIRCUITS...

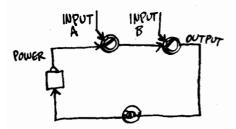




WHEN HO CURRENT PLOWS
THROUGH THE INPUT WIRE.
THE SWITCH REMAINS
OPEN, AS PICTURED
ABOVE. WHEN AN INPUT
SIGNAL ARRIVES, HOWEVER,
THE ELECTRONIC EQUIVALENT
OF A MINIATURE BOXING
GLOVE "PUNCHES" THE
SWITCH CLOSED,
RESULTING IN AN
OUTPUT SIGNAL.



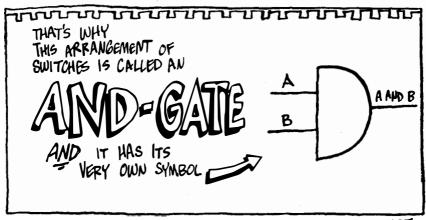
WHAT IS THE OUTPUT WHEN TWO SWITCHES (A, B) ARE ARRANGED IN SERIES, ONE APTER THE OTHER? [IN OUR DIAGRAM, PLEASE NOTE THE REARRANGEMENT OF WIRES, MADE FOR CONVENIENCE OF ILLUSTRATION.]



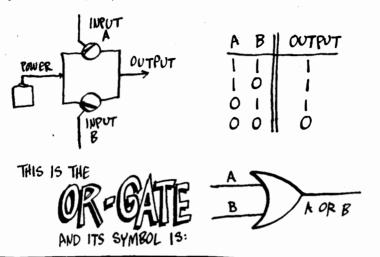
THE CURRENT CAN
FLOW ONLY IF BOTH
SWITCHES ARE CLOSED—
I.E., WHEN INPUT SIGNALS
ARRIVE SIMULTANEOUSLY
AT A AND B.

WRITING L FOR CURRENT AND O FOR NO CUFRENT, WE CAN THEN WRITE THIS IMPUT OUTPUT TABLE. LOOK FAMILIAR? IT SHOULD! IT'S IDENTICAL TO THE TRUTH TABLE FOR AND!

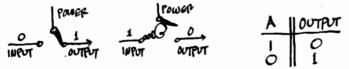
A	B	OUTPUT
T	1	1
1	0	0
0	1	0
0	0	



TWO SWITCHES CONNECTED IN PARALLEL BEHAVE LIKE LOGICAL OB: CURRENT CAN PASS FROM POWER TO OUTPUT IF EITHER SWITCH A, B IS CLOSED (OR IF BOTH ARE).



NOT IS NOT ANY MORE DIFFICULT... IT USES A SPECIAL SWITCH THAT REMAINS CLOSED UNTIL AN INPUT SIGNAL OPENS IT — JUST THE REVERSE OF AN ORDINARY SWITCH:



THIS KIND OF SWITCH IS CALLED

NOT-A

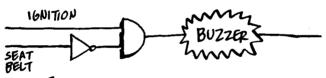
AND IT HAS A SYMBOL, TOO:

AN EVERYDAY EXAMPLE SHOWS HOW THESE SIMPLE GATES CAN MAKE LOGICAL DECISIONS.

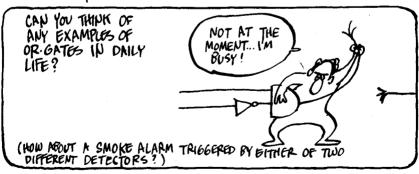
YOU KNOW
THOSE BUZZERS THAT
GO OFF WHEN YOU
START YOUR CAR
AND YOUR SEAT
BELT ISN'T FASTENED?
THE KIND THAT'S
SPECIALLY DESIGNED
TO PENETRATE
HUMAN BONE?



WELL, THAT'S BECAUSE THE SEAT BELT AND IGNITION ARE CONNECTED BY AN AND-GATE! LIKE SO:

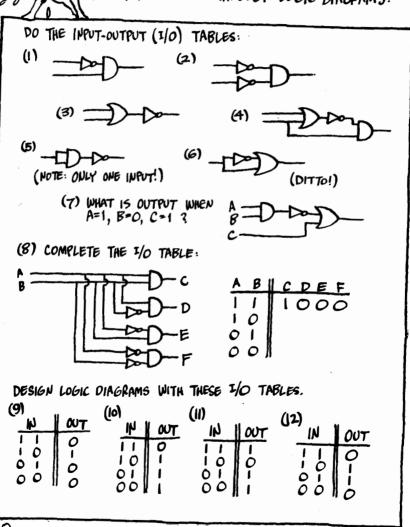


THAT IS, IF THE IGHTION IS ON AND THE SEAT BELT IS NOT, THE BUZZER SOUNDS! PRETTY LOGICAL, NO?

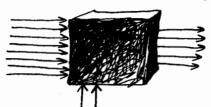




HERE ARE A PEW WARM-UP EXERCISES FOR CHASING THROUGH LOGIC DIAGRAMS:

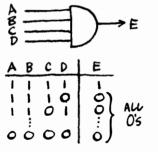


LOGIC GATES HAVE ONLY ONE OR TWO INPUTS AND A SINGLE OUTPUT — BUT COMPUTER COMPONENTS HAVE MANY INPUTS AND OUTPUTS WITH COMPLICATED INPUT/OUTPUT BEHAVIOR:



THE WONDERPUL FACT IS THAT ANY INPUT OUTPUT TABLE CAN BE PRODUCED BY A COMBINATION OF LOGIC GATES!

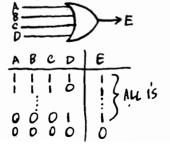
TO DO IT, YOU NEED MULTIPLE-INPUT LOGIC GATES. HERE'S A 4-INPUT **AND**-GATE:



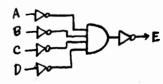
THIS MEANS
E=1 IF A=B=C=D=1
AND E=0 OTHERWISE?
THE GATE CAN BE MADE
WITH FOUR SWITCHES IN
SERIES:

A B C D

SIMILARLY, THERE'S A MULTIPLE-INPUT OR GATE:



IT CAN ACTUALLY BE MADE FROM AN AND-GATE AND SOME INVERTERS:



AS AN EXAMPLE OF HOW TO PRODUCE A GIVEN INPUT/OUTPUT TABLE, LET'S SOLVE PROBLEM #12:

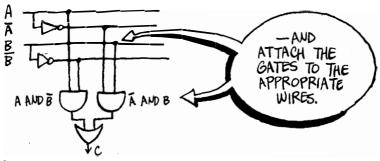


THE TABLE SAYS C=1 IF A=1 AND B=0 OR A=0 AND B=1. C=0 OTHERWISE.

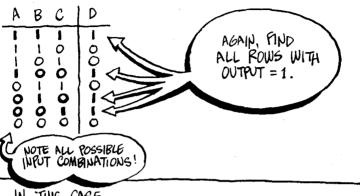
WRITING \overline{A} FOR NOT-A, This amounts to saying C=1 If A=1 AND $\overline{B}=1$ OR $\overline{A}=1$ AND B=1. C=0 OTHERWISE.

IN OTHER WORDS, $C=(A \text{ AND } \overline{B})$ OR $(\overline{A} \text{ AND } B)$

TO DRAW THE GROUT, RUN THE INPUT WIRES AND THEIR NEGATIVES IN ONE DIRECTION —



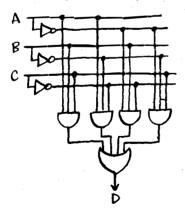
EXACTLY THE SAME METHOD WORKS FOR MORE INPUTS. FOR EXAMPLE:

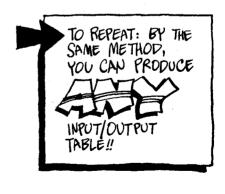


IN THIS CASE,

D = (A AND B AND C) OR (A AND B AND C) OR (A AND B AND C).

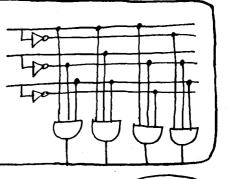
RUN THE INPUTS AND THEIR NEGATIVES ACROSS THE PAGE, ATTACH AND-GATES, THEN RUN THEM THROUGH AN OR-GATE!

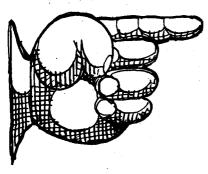






BY HOW YOU MAY
BE GETTING THE IDEA
THAT INFORMATION IS
ENCODED INSIDE COMPUTERS
AS STRINGS OF 1'S
AND 0'S, WHICH CAN BE
TRANSFORMED IN ANY
WAY WE LIKE BY THE
RIGHT COMBINATION OF
LOGIC GATES.





BUT WE NAVEN'T REPOLLY SEEN HOW LOGIC GATES CAN DO THE JOB COMPUTERS WERE DESIGNED FOR:



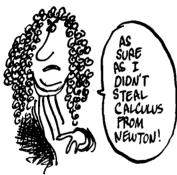


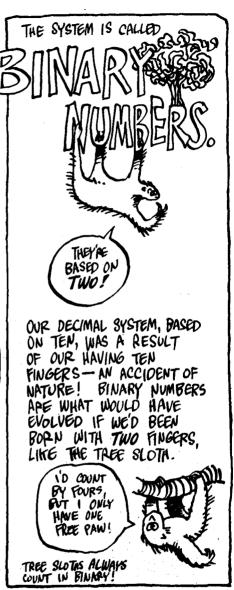
The questions:

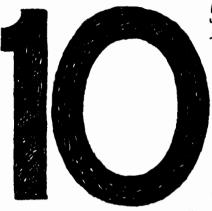
OS THERE SOME
NATURAL WAY
TO REPRESENT
NUMBERS USING
ONLY O'S AND I'S?
CAN THE OPERATIONS
OF ARITHMETIC
BE BUILT OUT OF
LOGIC?

The answer

(WHICH GOES BACK TO OUR OLD PAL LEIBNIZ):







LOOK AT THE SYMBOL

"10" — "ONE-ZERO." FORGET
THAT IT USUALLY MEANS
TEN! FORGET IT! STOP
CALLING IT THAT! IS
THERE ANYTHING THERE
THAT SAYS "TEN?"

NO!! IT'S JUST A ONE
FOLLOWED BY A ZERO —
IN AND OF ITSELF, IT
HAS NOTHING TO DO
WITH TEN!!!

THE SYMBOL DULY MAKES "TEN" FLASH THROUGH YOUR MIND BECAUSE YOU'VE ALWAYS CALLED IT THAT... IT'S LIKE A RITUAL: PERFORM IT OVER AND OVER AND IT BECOMES AUTOMATIC!



IN ACTUALITY, "10" MEANTS:

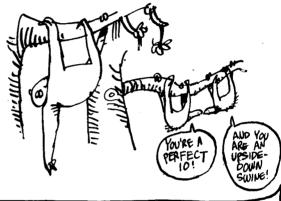


*REMEMBER - ON P. 24, WE AGREED TO CALL TEN FINGERS, NOT FINE, A HUMAN HANDFUL!



SINCE WE HUMANS HAVE TEN MINGERS, OUR "10" IS TEN... BUT TO AN ORGANISM WITH, SAY, FIGHT PINGERS, 10 WOULD MEAN EIGHT!





SOWE CAIN WRITE'S

10 BINARY = 2 DECIMAL

NOTE: DO NOT READ THIS AS "TEN EQUALS TWO."

TEN ODES NOT EQUAL TWO!! "ONE-ZERO IN BINARY"

EQUALS TWO!!





LIKEMSE, 100 — "ONE-ZERO-ZERO" — MEANS

1 HAMDFUL OF HAMDFULS.

IN DECIMAL, THAT'S 10×10, OF A HUNDRED. WELL, IN BINARY IT'S 10×10 ALSO—BUT THAT OMY AMOUNTS TO FOUR!

1000 15

10x 10x0 = 2x2x2 = 8 : and generally.

1 FOLLOWED BY N ZEROES K:

2 x ··· x 2 = 2 N

("TWO TO THE NTH POWER").

IN THE COMPUTER AGE, EVERYOUF WILL BE REQUIRED BY LAW TO MEMORIZE THE POWERS OF TWO, UP TO 2". BETTER NOT WAIT! AVOID JANK AND DO IT NOW!



ALL OTHER BINARY
NUMBERS — 101, 1111,
11000, AND EVERY OTHER
PATTERN OF 0'S AND
1'S — 15 A SUM OF SUCH
POWERS OF TWO!
1T'S COMPLETELY ANALO6005 TO DECIMAL.

N DECIMAL:	L IN BINARY,
497 =	111110001 =
400 + 90 + 7	100000000 256 + 10000000 128 + 1000000 64 + 100000 32 + 10000 16 + 10000 16

TO TRANSLATE A BINARY NUMBER INTO THE DECIMAL SYSTEM, LIST THE POWERS OF TWO OVER THE CORRESPONDING PLACES, AND ADD THOSE LYING OVER A 1.

$$\frac{2^{\circ} 2^{\circ} 2^$$

NOW YOU DO IT. CONVERT TO DEGIMAL:

(1) 11 (2) 101 (3) 1111111 (4) 11010101011101

TO MAKE THIS A BIT MORE CONCRETE -BINAPY DECIMA HERE'S HOW TO COUNT UP FROM 1 IN BINARY IT'S JUST LIKE COUNTING IN DECIMAL, ONLY EASIER. IN DECIMAL. TO COUNT PAST A 9, YOU WRITE O AND CARRY 1. IN BINARY. ġ YOU HAVE TO CARRY 1 EVERY OTHER NUMBER 1 12. ETC! | ETC! AS YOU MAY HAVE NOTICED, BINARY NUMBERS VERY FAST! THIS MAKES THEM HARD FOR US HUMANS TO USE WITHOUT MAKING MISTAKES -BUT FOR COMPUTERS THEY'RE IDEAL !

BINARY CALCULATION IS SIMPLE. THERE ARE ONLY PIVE RULES TO REMEMBER:

$$0 + 0 = 0$$

$$0+1=1$$

$$1 + 0 = 1$$

AND THE HANDY PIFTH RULE:



AS OPPOSED TO 100 SUMS IN DECIMAL: 9+6, 7+5, 9+3, 8+4, 4+6, etc etc Etc!!!

TO ADD TWO BINARY NUMBERS, PROCEED PLACE BY PLACE FROM RIGHT TO LEFT, CARRYING A 1 WHEN NECESSARY. HERE'S A STEP-BY-STEP EXAMPLE:

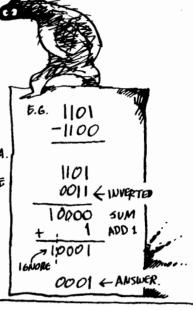
A FEW SUMS TO PRACTICE ON:





SUBTRACTION IS DONE BY ADDING !!

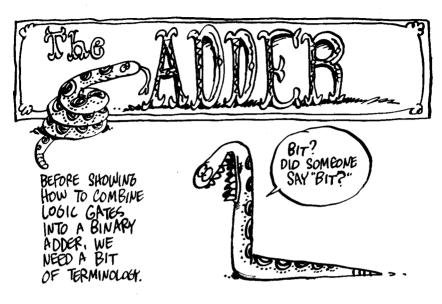
THE METHOD IS CALLED USING "TWO'S COMPLEMENT." FIRST YOU INVERT THE NUMBER TO BE SUBTRACTED, SO THAT ALL ITS I'S BECOME O'S AND VICE VERSA. THEN ADD THE TWO NUMBERS AND ADD I TO THE SUM. IGNORE THE PINAL CARRY AND THAT'S THE ANSWER!

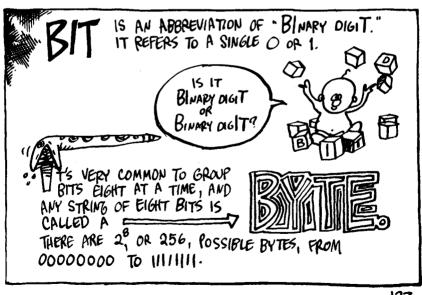


BINARY MULTIPLICATION—AND ANY MULTIPLICATION—MAY ALSO BE DONE BY REPEATED ADDITION: TO MULTIPLY A × B, JUST ADD A TO IT:SELF B TIMES. LIKEWISE, DIVISION CAN BE DONE BY REPEATED SUBTRACTION.

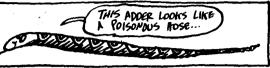
110 × 11 = 110 +110 +110 10010

The computer can do all arithmetic by adding !!



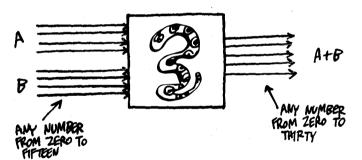


NOW LET'S SEE WHAT AN ADDER MIGHT LOOK LIKE.

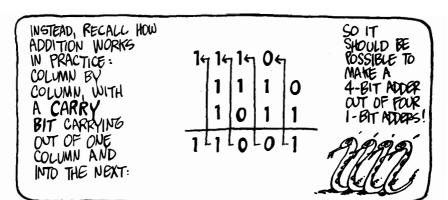


TO SAVE DRAWING, WE'LL MAKE IT A FOUR BIT ADDER, CAPABLE OF ADDING TWO 4-BIT NUMBERS, OR "NIBBLES." (YES, THEY'RE REALLY CALLED THAT!)

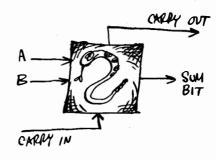
A= 1110 B= 1011 11001 THE INPUT OF OUR ACCER MUST CONSIST OF EIGHT BITS, FOUR FOR EACH NIBBLE. THE OUTPUT MUST BE FIVE BITS, THAT IS, A NIBBLE PLUS ONE BIT FOR A POSSIBLE CARRY. LIKE 80:

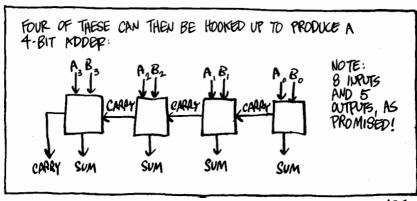


HOW TO PROCEED? ONE WAY IS TO MAKE A GIANT TRUTH TABLE, MATCHING EVERY POSSIBLE COMBINATION OF INPUTS WITH THE CORRECT OUTPUT, AND CONSTRUCTING A HUGE STEW OF ANDS AND NOTS TO FORCE A SOLUTION. THIS IS POSSIBLE, BUT THE COMPLEXITY OF THE TASK MIGHT OR JUST MAKE YOU THROW UP YOUR HANDS.



THE 1-BIT ADDER MUST HAVE THREE INPUTS — ONE FOR EACH OF THE TWO SUMMAND BITS OND ONE FOR THE BIT CARPIED IN — AND TWO OUTPUTS — ONE SUM BIT AND ONE CARRY-OUT BIT.

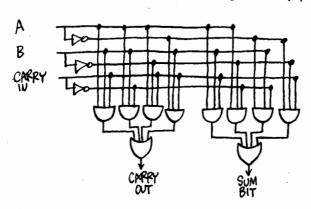




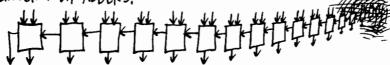
THE IN	PUT/OUTPUT FOR THE
TABLE	FOR THE
1-BIT	ADDER:
-	• ,

A	В	CARRY	CARRY SUM OUT BIT
1 1 0 0 0 0	1-0000	-0-0-0-0	0-0-0

NOW THERE'S NOTHING TO IT! REMEMBER, LOGIC GATES CAN BE RIGGED UP TO PRODUCE ANY INPUT/OUTPUT TABLE. IN THIS CASE, JUST TREAT EACH OUTPUT COLUMN SEPARATELY:



YOU CAN ADD TWO NUMBERS OF ANY LENGTH BY HOOKING TOGETHER ENOUGH 1-BIT ADDERS.





THE IMPLICATION OF THE LAST TWO SECTIONS IS THAT BINDRY IS THE "NATURAL" SYSTEM FOR ENCODING NUMBERS IN A MACHINE MADE OF ON/OFF SWITCHES. EVEN SO, COMPUTERS USE SEVERAL VARIATIONS ON THE BASIC IDEA.

INTEGERS, OR WHOLE NUMBERS—IF THEY AREN'T TOO LARGE—ARE ENCODED IN STRAIGHT BINARY. FOR INSTANCE,

WOULD BECOME

10111001

FLOATING POINT REPRESENTATION IS FOR LARGE OR PRACTIONAL NUMBERS. FOR EXAMPLE, 19,700,030.2 WOULD BE ENCODED AS THE BINMY EQUIVALENT OF 197 5

MEANING 197 × 105.
FLOATING POINT REPRESENTATION OFFEN INVOLVES ROUNDING OFF.

BINARY CODED DECIMAL
REPRESENTS A NUMBER IN
DECIMAL, BUT WITH EACH DIGHT
ENCODED IN BINARY. 967,
FOR INSTANCE, WOULD BECOME

001 0110 0111

AND WHAT ABOUT NON-NUMERICAL INFORMATION — THE ALPHABET, PUNCTUATION MARKS, OTHER SYMBOLS, AND EVEN THE BLANK SPACE ??

NEXT

FOUR

BITS

SINCE THERE IS NO NATURAL WAY
TO ENCODE THESE
INTO O'S AND 1'S,
COMPUTER SCIENTISTS
INVENTED AND
ADOPTED A
STANDARD CODE
BY MUTUAL
AGREEMENT:

ASCII

THE AMERICAN STANDARD CODE FOR INFORMATION INTERCHANCE

(ACTUALLY, ASCIL IS USED BY EVERYONE BUT IBM, WHICH HAS ITS OWN CODE, CALLED EBCDIC.)

ONTY!
WHERE
IS
UNITY!

FIRST THREE BITS

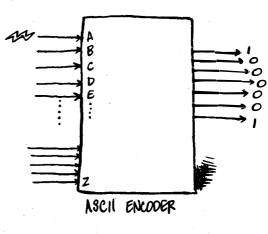
		000	00-	0-0	0-1	100	-0-	1-0	
1	0000	NUL	DLE	SP	0	@	Р	,	P
ı	0001	SOH	DC1	!	1	Α	Q	а	9
	0010	Stx	002	11	2	₽	R	ط	r
	0011	etx	003	#	3	С	S	c	5
	0100	EOT	DC4	\$	4	D	٢	d	t
	0101	enq	NAK	%	5	E	υ	e	u
ı	0110	ack	SYV	&	6	F	٧	£	v
	0111	BEL	ET8	′	7	G	W	9	w
	1000	BS	CPN	(8	Н	Х	h	х
	1001	HT	EM)	9	I	Y	i	y
	1010	LF	SUB	*	:	J	Z	j	z
	1011	VT	ESC	+.	;	K	Ľ	k	{
	1100	FF	۴S	,	<	L	\	١	1
	1101	CR	65	-	=	М]	m	}
1	1110	SO.	RS		>	N	Λ	n	~
-	1111	S١	US	/	?	0	_	0	DEL

THUS, THE LETTER "T" IS ENCODED AS

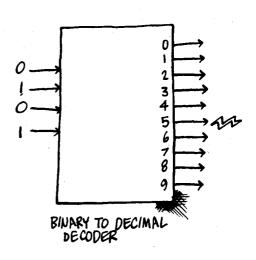
ATTHE FIRST TWO COLUMNS CONTMN SYMBOLS FOR SUCH THINGS AS "START OF HEADING" (SOH) AND OTHER TEXTUAL DIRECTIONS.

TO ENCODE AND DECODE DATA, COMPUTERS USE LOGIC DEVICES CALLED, NATURALLY ENOUGH, **ENCODERS** AND **DECODERS**.

AN ENCODER
USUALLY HAS MANY
INPUTS AND A
FON OUTPUTS. A
SINGLE INPUT SIGNAL
PRODUCES A
PATTERN OF
OUTPUTS. FOR
EXAMPLE, A
COMPUTER KEYBOARD
IS ATTACHED TO AN
ENCODER INHICH
TRANSLATES A
SINGLE KEYSTROKE
INTO ITS ASCII
CODE.

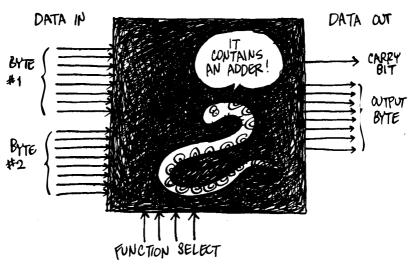


A CLECULED WAY AROUND, TRANSLATING A PATTERN OF BITS INTO A SINGLE OUTPUT SIGNAL. ONE DECODER CONVERTS A BINARY NIBBLE INTO A DECIMAL DIGIT. ANOTHER TRANSFORMS A SPECIAL LOCATION, OF ADDRESS, IN MEMORY INTO A SIGNAL TO THAT MEMORY CELL. (SEE P. 155.)

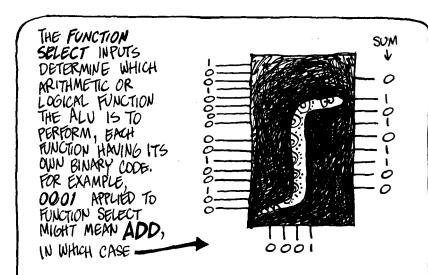


ONCE ALPHANUMERIC INFORMATION IS ENCODED IN BINARY STRINGS, IT IS READY TO BE PROCESSED BY THE CONPUTER'S MOST ELABORATE COMBINATION OF LOGIC BATES, THE

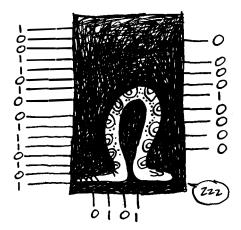
ARTHMETIC LOGIC UNIT (OR ALU, FOR SHORT),



THIS IS THE MACHINE'S CENTRAL PROCESSOR, WHICH CAN ADD, SUBTRACT, MULTIPLY, COMPARE, SHIFT, AND PERFORM A WEALTH OF OTHER LOGICAL FUNCTIONS. THE DRAWING ABOVE REPRESENTS AN B-BIT ALU, BUT THEY CAN RANGE FROM FOUR TO SIXTY BIT CAPABILITY, DEPENDING ON THE COMPUTER.



ANOTHER FUNCTION (0101, SAY) MIGHT COMPARE
TWO BYTES, BIT BY BIT, AND OUTPUT A 1 WHEREVER THEY AGREE.
(MEANWHILE, THE ADDER TAKES A UMP.)



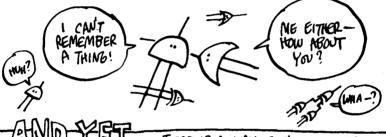
YOU CAN GET AN IDEA OF A FANCY ALU'S CAPABILITIES FROM THE LIST ON PAGE 182.



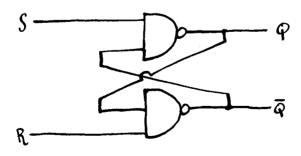




VERSATILE AS THEY MAY BE, THE LOGICAL COMBINATIONS WE'VE BEEN SKETCHING STILL HAVE NO MEMORY. THEIR OUTPUT CONTINUES ONLY AS LONG AS THE INPUT IS APPLIED.



INDIFFE IS A WAY TO HOOK THESE LOGICAL BUT SEMILE GATES TOGETHER INTO A GADGET THAT HOLDS AN OUTPUT INDEFINITELY: THE FLIP-FLOP. STARE AT THIS A MINUTE!



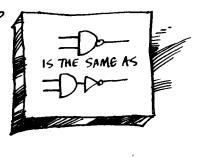


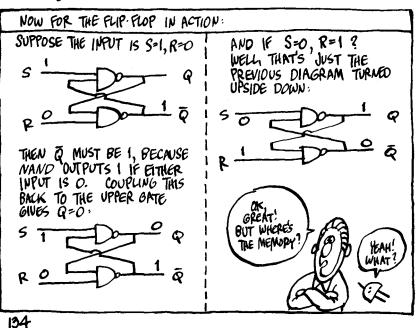
BESIDES THE STRANGE WAY A FLIP-FLOP EATS ITS OWN TML, PLEASE NOTE THE UNFAMILIAR GATE USED IN THE CONSTRUCTION. IT'S CALLED A

NANDO

WHICH IS MERELY AN ABBREVIATION OF "NOT AND."

Pr	8	NAN	0
1-00	-0-0	0	

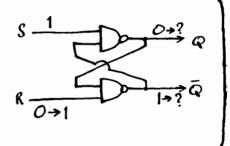




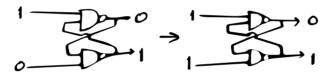
NOW WHAT HAPPENS
WHEN THE INPUT

CITY AND S

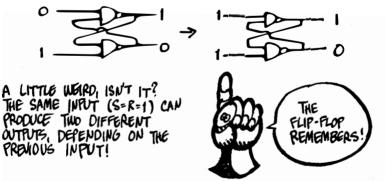
SUPPOSING WE BEGIN WITH
THE INPUT (S=1, R=0),
WHAT DOES CHANGING
IT TO (S=1, R=1)
DO TO THE FUP-FLOP'S
OUTPUT?



THE ANSWER IS: INCUTATIONS THE LOWER NAND-GATE'S INPUT BECOMES (0, 1), SO ITS OUTPUT Q IS STILL 1, SO Q REMAINS O.



But precisely the SAME line of Reasoning Shows no change in output when input changes to (s=1, r=1) from (s=0, R=1):

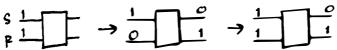


THE WAY A FLIP-FLOP IS USED IS THIS: IT BEGINS BY SITTING THERE WITH A CONSTANT INPUT OF (S=1, R=1) AND AN OUTPUT OF GOD-KNOWS-WHAT:



YOU SET THE FLIP-FLOP [I.E., MAKE Q=1] BY FLASHING A O MOMENTARILY DOWN THE S-WIRE, AND THEN RETURNING IT TO 1:

OF YOU CAN RESET IT [MAKE Q=0] BY FLASHING A O DOWN THE R.WIRE, THEN RETURNING IT TO 1:



IN ETHER CASE, AS LONG AS
(1,1) KEEPS
COMING IN, THE
FLIP-FLOP WILL
MAINTAIN ITS
OUTPUT UNTIL
IT'S CHANGED WITH
ANOTHER INCOMING
O.



THE ONLY INPUT COMBINATION WE HAVEN'T CHECKED IS (R=S=0). It's EASY TO VERIFY THAT IT PRODUCES OUTPUT OF Q=Q=1:



WHAT HAPPENS WHEN THE INPUT RETURNS TO (1, 1)?



THE ANSWER IS NOT SO CLEAR: IT DEPENDS ON WHICH OUTPUT HAPPENS TO FLOP FIRST!! (ONE OF THEM MUST.)



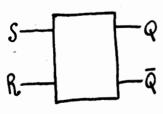
HOWEVER:

SINCE THERE IS IND IWAY OF KNOWING WHICH OF THESE WILL ACTUALLY HAPPEN, AND WE DON'T WANT OUR FUR-FLOPS IN RANDOM STATES, THE INPUT (5=0, R=0) IS



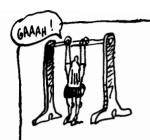
DISALLOWED.

WE CAN SUMMARIZE THE BASIC "R-S" FLIP-FLOP LIKE SO:



S	R	QQ
T	10	NO CHANGE
0	0	DISALLOWED!

FLIP-FLOP INPUTS ARE ALWAYS ARRANGED TO MAKE CERTAIN THE DISALLOWED STATE CANNOT ARRIVE.



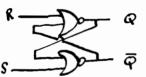
A LITTLE EXERCISE:



_	/			
	١	AM	1	
	Ť	HE	1	1
	4	NT US	1	い
	11	אן טי	۱:/	$\neg \! \prime$
•	•			

A	В	11 NOR_
<u> 0<</u>	3-0-	000

A BASIC RS FUP. FLOP MAY ALSO BE MADE OUT OF NOR-GATES:





- 1. WHAT IS THE OUTPUT WHEN R=0, S=1 ? WHEN S=0, R=1 ?
- 2. WHAT HAPPENS WHEN EACH OF THESE INPUT CONDITIONS CHANGES TO R=0, S=0?
- 3. WHAT IS THE OUTPUT WHEN R=1, S=1? WHAT HAPPENS WHEN THIS CHANGES TO R=0, S=0?
- 4. WHAT IMPUT COMBINATION MUST BE DISALLOWED?
- 5. IF R=0, S=0, How Do YN SET THIS PUP: FLOP (I.E., MAKE Q=1)? HOW DO YOU RESET IT?

By the way, a flip-flop is also called a LATCH, because it "Locks in" data.



BLEENCH!

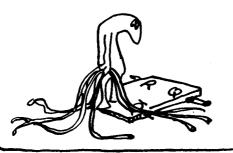
REGISTIERS, CLITCHES & STATES

IF THE FLIP-FLOP IS A DEVICE FOR STORING ONE BIT, A REBOISTER STORES SEVERAL BITS SIMULTANEOUSLY. IT'S LIKE A ROW OF BOXES, EACH HOLDING ONE BIT.

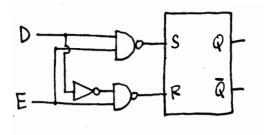


A ROW OF FLIP-FLOPS SHOULD DO THE JOB :...

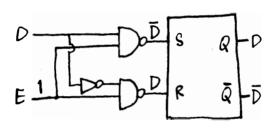




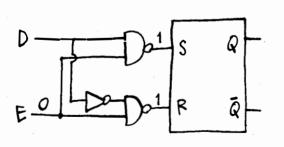
... SORT OF! BOP IF YOU TRY AND MAKE THIS WORK BY HOOKING UP SOME INPUTS TO RS FLY-FLOPS, YOU MAY PIND YOURSELF GROWING CONFUSED! THE SOLUTION IS TO ADD A "GATING NETWORK" TO THE BASIC R-S FLIP-FLOP.



HERE "D" STANDS POR DATA, AND
"E" STANDS FOR ENABLE. NOTE
THAT THE GATING NETWORK MAKES
IT IMPOSSIBLE FOR R AND S TO BE
ZERO SIMULTANEOUSLY.

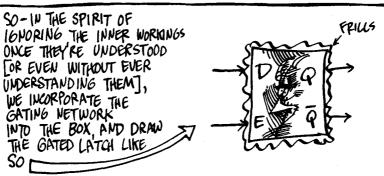


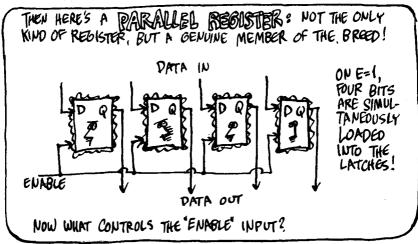
WHEN E=1, THEN
R=D AND S=D
(NOT-D). HENCE,
THE VALUE OF D
IS STORED AT Q.
IN OTHER WORDS,
E=1 ENABLES
THE BIT D TO
BE LONDED INTO
THE FLIP-FLOP.

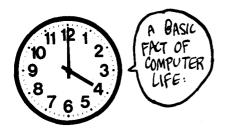


WHEN E=O,
S AND R BOTH
BECOME 1, AND
THE FUP-FLOP
DOES NOT CHANGE.
THAT IS, E=O
BLOCKS THE
APRIVAL OF MORE
DATA.









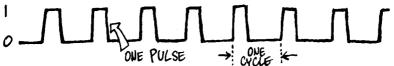
AS SOON AS YOU BEGIN
STORING DATA, QUESTIONS
OF TIMING ARISE: HOW LONG
DO YOU STORE IT? WHEN
DO YOU MOVE IT? HOW DO
YOU SYNCHRONIZE SIGNALS?
THESE ISSUES ARE SO CRITICAL
THAT LOGIC WITH MEMORY

IS CALLED SECUENTIAL, TO DISTINGUISH IT FROM THE PURELY COMBINATIONAL LOGIC OF MEMORY-LESS NETWORKS. TO KEEP THE SEQUENTIAL LOGIC IN STEP,

ALL COMPUTERS HAVE CLOCKS!

THE CLOCK'S PULSE IS THE COMPUTER'S HEARTBEAT—ONLY INSTEAD OF A WARM, RABGED HUMAN HEARTBEAT, LIKE THIS—

THE COMPUTER'S PULSE IS SQUARE AND COLD:



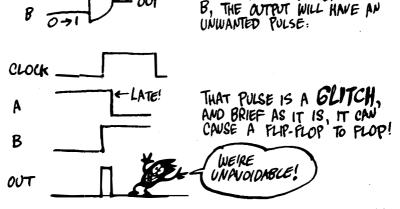
ONE CLOCK PULSE IS THE BURST OF CURRENT WHEN CLOCK OUTPUT = 1. ONE CYCLE IS THE INTERVAL FROM THE BEGINNING OF A PULSE TO THE BEGINNING OF THE NEXT. DEPENDING ON THE CONTUTER, THE CLOCK FREQUENCY MAY BE HUNDREDS OF THOUSANDS TO BILLIONS OF CYCLES PER SECOND!

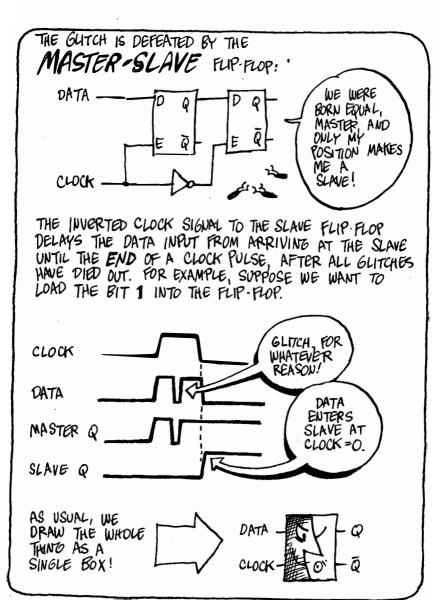
Slow Compotep:		+ 100000 SEC	
-------------------	--	--------------	--

THE IDEA OF USING A CLOCK IS THAT THE COMPUTER'S LOGICAL STATE SHOULD CHANGE GOODY ON THE CLOCK PULSE. IDEALLY, WHEN THE CLOCK HITS 1, ALL SIGNALS MOVE, THEN STOP ON CLOCK = O. THEN GO ... THEN STOP ... THEN GO ... STOR GO STOR ... A TYPICAL EXAMPLE THEN A NEW IS TO ATTACH THE bit of CLOCK TO THE "ENABLE" DATA IS IMPUT OF A GATED LATCH, LOADED AT IN WHICH CASE THE LATCH every BECOMES KNOWN AS A CLOCK "D FLIP-FLOP." PULSE! CLOCK

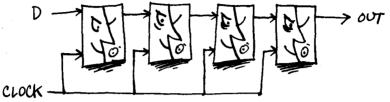
UNFORTUNATELY, THINGS ARE RAPELY IDEAL! IT TAKES A NON-ZERO TIME FOR A SIGNAL TO PASS ALONG A WIRE, SO THINGS ARE PEUR PERFECTLY SYNCHRONIZED. FOR EXAMPLE, SUPPOSE AT AN AND GATE, ONE IMPUT IS CHANGING FROM 1 TO 0, AND THE OTHER FROM 0 TO 1:

IF A CHANGES AFTER





STRINGING A NUMBER OF MASTER-SLAVE FLIP-FLOPS
TOGETHER MAKES A SINFT RECOSTERS

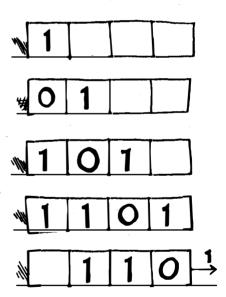


DATA ENTER A SHIFT REGISTER ONE BIT AT A TIME, SHIFTING TO THE RIGHT WITH EACH NEW CLOCK PULSE.

FOR EXAMPLE, THE NIBBLE 1101 WOULD ENTER THE SHIFT REGISTER LIKE THIS:

EACH CLOCK PULSE
BRINGS A NEW BIT
INTO THE REGISTER.
(WHY DOESN'T THE
BIT TRAVEL ALL THE
WAY THROUGH ON
ONE PULSE? BECAUSE
OF THE MASTER-SLAVE
FLIP-FLOPS!)

LIKAWISE, THE NIBBLE SHIFTS OUT ONE BIT AT A TIME.



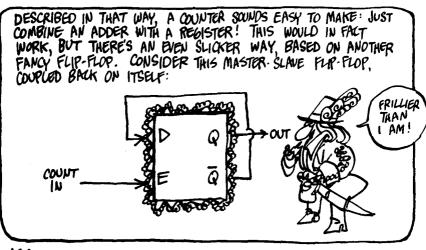


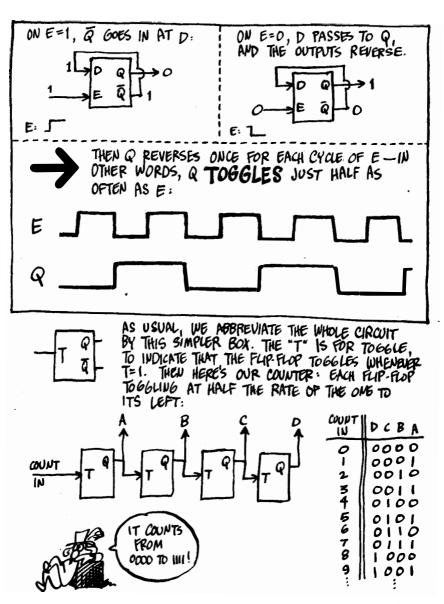
SHIFT REGISTERS ARE USEFUL WHEN INFORMATION IS TO BE TRANSMITTED SERIALLY, OR ONE BIT AT A TIME.



A COUNTER IS JUST WHAT
IT SOUNDS LIKE: SOMETHING
THAT COUNTS. IN OTHER
WORDS, IT'S A REGISTER
THAT INCREMENTS ITSELF—
ADDS 1 TO ITS CONTENTS—
WHENEVER A "COUNT"
SIGNAL ARRIVES:

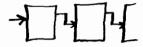
ETC!





A FEW ITEMS OF NOTE:

THIS COUNTER IS CALLED AN "ASYNCHRONOUS RIPPLE COUNTER," BECAUSE THE COUNT RIPPLES THROUGH FROM ONE FLIP-FLOP TO THE NEXT. THIS CAUSES A SUGHT DELAY BEFORE THE COUNT IS REGISTERED.





When the 16th Count Pulse Arrives, The Counter Returns To O. to 60 HIGHER THAN 15, MORE FLIP-FLOPS ARE NEEDED.

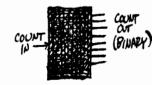
THIS 14. BIT COUNTER CAN GO FROM O TO 214-1 = 16.383

NITERUAL C

THE NTH FUP-FLOP IN A RIPPLE COUNTER **DIVIDES** THE INCOMING PULSE BY 2°. THIS IS THE PRINCIPLE ON WHICH DIGITAL WATCHES ARE BASED: A HIGH-FREQUENCY INTERNAL CLOCK PULSE IS DIVIDED TO A RATE OF PREUSELY ONE CYCLE RER SECOND.



THERE ARE ALSO SYNCHRONOUS COUNTERS, WHICH REGISTER ALL BITS SIMULTANEOUSLY, AND COUNTERS WHICH RETURN TO O ON ANY PAEASSIGNED NUMBER. IN ANY CASE, FROM NOW ON, A COUNTER IS JUST ANOTHER BLACK BOX!!





THE AMAZING NAND:



1. SHOW THAT



A DOLD IS THE AS ID-



CONCLUDE THAT => ALL LOGIC CAN BE DERIVED FROM THE SINGLE RELATION NAME!!

1. CAN THE SAME BE DONE WITH NOR?

3. SHOW THAT



IS THE SAME AS



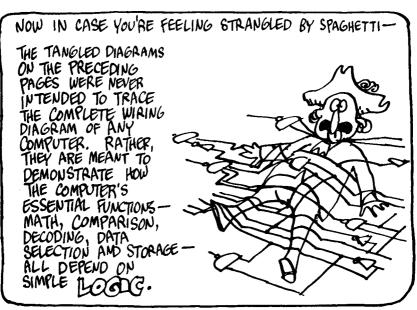
REDRAW THE ADDER ON P. 126 USING ONLY NAND-GATES.

4. GIVEN A 4-BIT SHIFT REGISTER,



Show its contents after each of four clock pulses as the NIBBLE DOIL IS ENTERED.

- 5. HOW WOULD YOU ATTACH
 A BUZZER TO A COUNTER
 TO SOUND WHEN THE
 COUNT HITS UINE
 (=1001 IN BINARY)?
 HINT: LOOK AT THE SEAT
 BELT BUZZER ON P. 109.
- G. CONVINCE YOURSELF THAT
 ATTACHING INVERTERS TO
 THE OUTPUTS MAKES A
 COUNTER COUNT BACKWARDS.



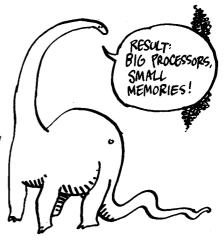
NOW THAT YOU PRESUMABLY BELIEVE IN THE POWER OF LOGIC, NO MORE WIRING DIAGRAMS ARE NEEDED!







N THE INFANCY
OF ELECTRONIC
COMPUTING, MEMORY
WAS ALWAYS MORE
EXPENSIVE THAN SHEER
COMPUTING POWER.
PLENTY OF PROCESSING
COULD BE DONE WITH
RELATINELY FEW COMPONENTS,
BUT EVERY INCREASE IN
MEMORY SIMPLY MEANT
MONES — MORE ACTUAL,
PHYSICAL PLACES TO STORE
THANGS!



SINCE THEN, RESEARCH INTO MEMORY TECHNOLOGY HAS BROUGHT DOWN THE COST CONSIDERABLY. FOR A FOW HUNDRED DOLLARS YOU CAN BUY A MICRO WITH OVER 64,000 BYTES OF MEMORY, COMPARED WITH BRIDGIE'S MEMORY OF ABOUT 100 NUMBERS*—AT A COST OF MILLIONS!!



*ENIAC DID NOT COMPUTE IN BINARY.



AN IMPORTANT DISTINCTION EXISTS RETWEEN

ELECTRONIC

AND

ELECTRO-MECHANICAL

MEMORY DEVICES.

ELECTRONIC MEMORIES, WITH NO MONING PARTS, ARE AS FAST AS THE REST OF THE COMPUTER.

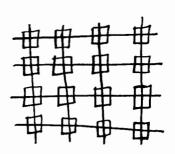
ELECTROMECHANICAL MEMORIES HAVE MOVING PAPTS, LIKE DISKS OR REELS OF TAPE. THIS MAKES THEM SLOW—HOW SLOW DEPENDING ON THE TYPE OF MEMORY.



ELECTRONIC MEMORIES' SPEED MAKES THEM IDEAL FOR THE COMPUTER'S MAIN, OR INTERNAL MEMORY, WHILE ELECTRO-MECHANICAL MEMORIES ARE USED FOR AUXILIARY STORAGE OUTSIDE THE MACHINE,

ELECTROMAGNETIC MEMORIES COMPENSATE FOR THEIR SLOWNESS WITH A GIGANTIC CAPACITY. ONE HARD DISK CAN STORE UP TO TEN MILLION BYTES, COMPARED WITH A TYPICAL MICRO'S MAIN MEMORY OF 65,536 (=216) BYTES.

INTERNAL MEMORY
CAN BE THOUGHT OF AS
A SIMPLE GRID, WITH A
CELL AT EACH INTERSECTION.
DEPENDING ON THE
COMPUTER, EACH CELL CAN
HOLD ONE BYTE, TWO
BYTES, OR MORE.



EVERY CELL HAS A UNIQUE ADDRESS, SPECIFYING WHERE IT SITS IN THE GRID.

ADDRESS OO10 0011

IN PRACTICE, THERE MAY BE
MANY SUCH GRIDS, IN WHICH
CASE THE ADDRESS SPECIFIES
THE GRID NUMBER, AS WELL
AS THE ROW AND COLUMN
WITHIN IT.

DO NOT CONFUSE A CELL'S MODRESS WITH
IT'S CONTENTS!!

WHAT IS THE MAXIMUM NUMBER OF CELLS THE COMPUTER CAN MODRESS? THIS DEPENDS ON THE LENGTH AND STRUCTURE OF THE COMPUTER'S "WORDS." FOR EXAMPLE, A 32-BIT MACHINE MAY INTERPRET THE FIRST 8 BITS AS AN INSTRUCTION...

B-DIT INSTRUCTION

TOGOTTO OT

24-BIT ADDRESS

... AND THE REMAINING 24 BITS AS AN ADDRESS.

In that case, addresses can be anything between

00000 O

 $1 = 2^{24} - 1$

GIVING 284 POSSIBLE MEMORY CELLS.



AN B.BIT MICRO, ON THE OTHER HAND, MIGHT PROCESS THREE BYTES IN SUCCESSION:

00110111

AN INSTRUCTION

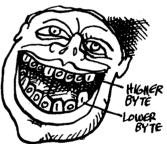
10011010

THE FIRST HALF OF AN ADDRESS,

00010100

AND THE SECOND HALF OF AN ADDRESS.

HERE THE ADDRESS
15 16 BITS LONG,
61VING 216-65,536
POSSIBLE ADDRESSES.



16-BIT WORDS ARE OPTEN SPLIT LIKE THIS INTO HIGHER-LEVEL AND LOWER. LEVEL BYTES.

10001101 00010010 HIGHER LOWER

156

TO MAKE ADDRESSES SHORTER AND MORE READABLE, THEY'RE OPTEN EXPRESSED IN

MEXADECIMAL, or base-16, numerals.

ETC!



JUST AS BASE 10 NUMBERS REQUIRE THE DIGITS O-9, SO HEXADECIMAL NEEDS DIGITS FROM O TO FIFTEEN. THE EXTRAS ARE REPRESENTED BY THE LETTERS A-F:

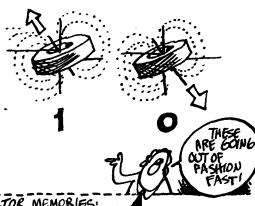
FOR EXAMPLE:

TO CONVERT BINARY TO HEX:
GROUP THE BINARY NUMBER INTO NIBBLES, STARTING FROM THE RIGHT. CONVERT EACH NIBBLE TO A HEX DIGIT!

TO CONVERT HEX TO BINARY, JUST REVERSE THE PROCESS.

FROM THE HARDWARE POINT OF VIEW, THERE ARE THREE MAIN TYPES OF INTERNAL MEMORY.

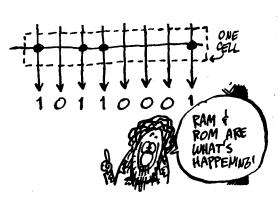


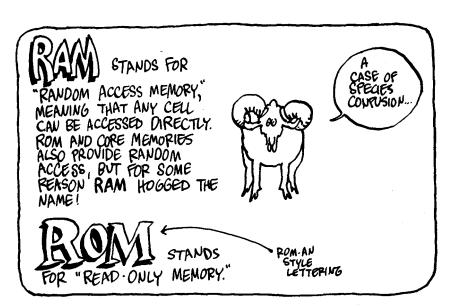


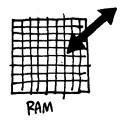
AND TWO SEMICONDUCTOR MEMORIES.

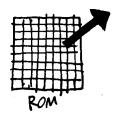
USES FLIP-FLOPS
TO STORE BITS—
SO ENCH MEMORY
CELL 15 ESSENTIALLY
A (PARALLEL) REGISTER!

INDICATES A 1 OR O AT EACH GRID POINT BY THE PRESENCE OR AN ELECTRIC CONNECTION THERE.









THE PRACTICAL DIFFERENCE BETWEEN THEM IS THAT YOU CAN ONLY READ WHAT'S IN ROM, WHILE WITH RAM YOU CAN READ THINGS OUT OR WRITE THEM IN WITH EQUAL EASE.

IN GENERAL!



when you load a program into the computer, it is stored in **RAM**. UNFORTUNATELY, RAM IS VOLATULE,

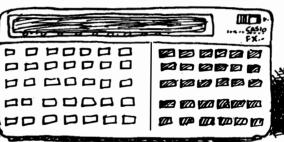




IT FORGETS EVERYTHING WHEN THE POWER IS TURNED OFF.

FOR EXAMPLE, I OWN A BATTERY-POWERED POCKET COMPUTER WITH 1680 BYTES OF RAM. IT CAN STORE UP TO TEN PROGRAMS EVEN WHEN I TORN IT OFF, BECAUSE IT KEEPS SOME ELECTRICITY RUNNING THROUGH MEMORY.

BUT WHEN THE BATTERY DIES... BYE BYE, PROGRAMS!



RAM VOLATILITY IS ONE REASON THAT THE MAGNIFICENT, INFALLIBLE COMPUTER IS VULNERABLE TO THE VAGARIES OF OUTMODED, ERRATIC POWER GENERATING STATIONS!





ROM — "READ-ONLY MEMORY"—
ONCE ITS CONTENTS ARE ENTERED,
CAN NEVER BE REWRITTEN.*
ORDINARILY ROM IS PROGRAMMED
AT THE FACTORY BUT THERE ARE
NOW ALSO PROMIS — PROGRAMMABLE

ROMS — WHICH CAN BE CUSTOM-PROGRAMMED TO THE USER'S

SPECIFICATIONS.

WHAT ARE YOU DOING ON PROM NIGHT!

GOING

SHACK

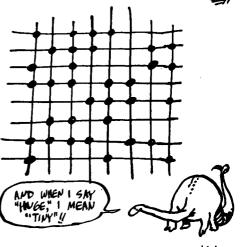
TO GET

PIKED.

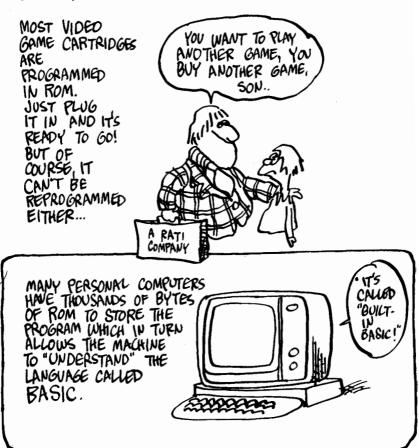
TO RADIO

*Except for EPROM — EPROABLE PROGRAMMABLE ROM'T DET UN'T DET IPTO THAT!

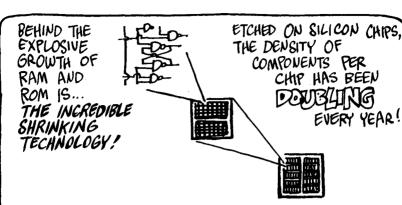
UNLIKE RAM, ROM
IS WON. VOLATIOLE:
IT KEEPS ITS
CONTENTS EVEN
WITHOUT POWER.
AFTER ALL, IT'S
NOTIKING BUT A
KUGE GRID OF WIRES
WITH PHYSICAL
CONVECTIONS AT SOME
INTERSECTIONS.
THE CONVECTIONS
REMAIN, REGARDLESS
OF ELECTRIC CURRENT.



SOME TYPICAL USES OF ROM:



AND, AS WE'LL SEE, ROM PLAYS AN IMPORTANT ROLE IN THE COMPUTER'S CONTROL SECTION.



THE STANDARD MEASURE OF CHIP STORAGE IS THE SHORT FOR "KILO" ("CHILO" IS GREEK FOR 1000), IN COMPUTERESE IT MEANS 2", THE POWER OF TWO CLOSEST TO 1000:



The Pirst RAM
CHIP WITH 1K
BITS OF STORAGE
WAS A
SENSATION—
BUT YOW
64 K IS
COMMON, AND THE
256K CHIP
HAS ARRIVED!
WHAT'S
NEXT?





THE ANSWER?

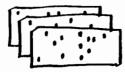
mass storage.

AS THE NAME IMPLIES,
MASS STORAGE IS MEMORY
THAT CAN STORE A LOT!!
ALMOST ALL MASS STORAGE
DEVICES ARE NON-VOLATILE
AND HAVE A MECHANICAL
COMPONENT THAT MAKES THEM
MUCH SLOWER THAN ELECTRONIC
RANDOM ACCESS MEMORIES.

DISCUM!



FOR EXAMPLE.



THE CARDS OF JACQUARD, BABBAGE, AND HOLLERITH ARE STILL IN USE!

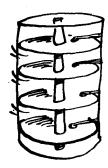
SAME IDEA AS PUNCHCARDS: A HOLE REPRESENTS 1, A NON-HOLE O.

STORES BITS AS SMALL MAGNETIC REGIONS, WHICH MAY BE MAGNETIZED IN ONE OF TWO DIRECTIONS, REPRESENTING 1 OR O.

FASTER, LESS BULKY, AND THE CURRENT STORAGE OF CHOICE IS THE



DISKS ALSO STORE BITS AS TIMY MAGNETIZED REBIONS — UP TO 10 MILLION BYTES PER DISK!



A BIG COMPUTER SYSTEM USUALLY WAS MULTIPLE DISK DRIVES, WITH PHONOGRAPH-ARMLIKE READ/WRITE HEADS DARTING BACK AND FORTH ACROSS THE WHIRLING PLATTERS.

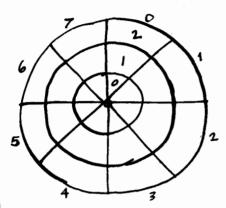


ARE SMALL,
LOW COST
MAGNETIC
DISKS MADE
OF PLASTIC.
THEY ALWAYS
STAY IN THER
JACKETS, BECAUSE
A SPECK OF
DUST CAN CREATE
A MONSTER
GLITCH!



OTHER, MORE EXOTIC MASS STORAGE TECHNOLOGIES INCLUDE BUBBLE MEMORIES, CHARGE-COUPLED DEVICES, AND OPTICAL DISK'S READ BY LASERS.

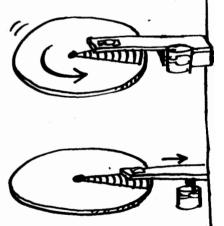
LIKE INTERNAL MEMORY, MASS STORAGE MUST BE ORGANIZED, OR "FORMATTED." TAKE THE FLOPPY DISK FOR EXAMPLE:



FLOPPIES ARE FORMATTED
INTO RINGS AND
SECTORS — THREE
RINGS AND EIGHT SECTORS,
IN THIS VERY OVERSIMPLIFIED DISK.
(IT'S MORE LIKE 26
SECTORS AND 77 RINGS
IN A GENUINE DISK.)

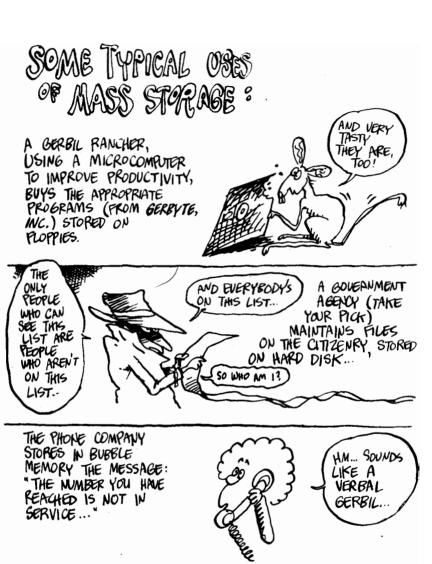
TO Access A PARTICULAR BLOCK OF DATA, YOU SPECIFY THE RING NUMBER AND SECTOR NUMBER. THEN THE DISK DRIVE

- 1) SPINS THE DISK UNTIL THAT SECTOR LIE'S UNDER THE READ/WRITE HEAD
- 2) MOVES THE HEAD IN OR OUT TO THE PROPER RING.





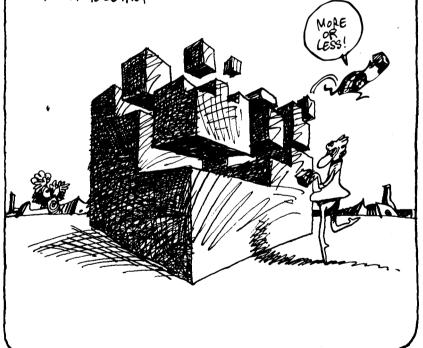
THIS PROCESS TAKES MILLISECONDS — AN ETERNITY TO A COMPUTER!

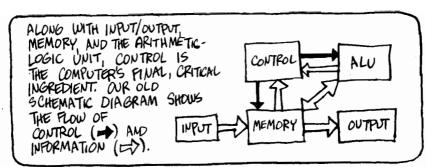


WELL, YOU GET THE PICTURE... NOW IT'S TIME TO MOVE ON...

GETTING EVERYTHING UNDER PIER

IN WHICH ALL THE BLACK BOXES ARE FINALLY SEEN TO FIT TOGETHER...

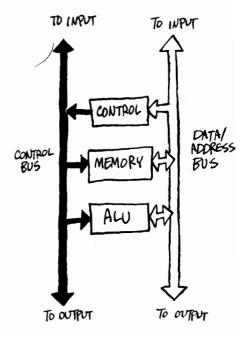




IT HELPS TO REDRAW THIS DIAGRAM IN A WAY THAT BETTER REFLECTS A GENUINE COMPUTER DESIGN KNOWN AS "BUS ARCHITECTURE."

THE VERTICAL ARROWS, REPRESENTING ELECTRICAL PATHWAYS A BYTE OR MORE WIDE, PRE THE BUSES.

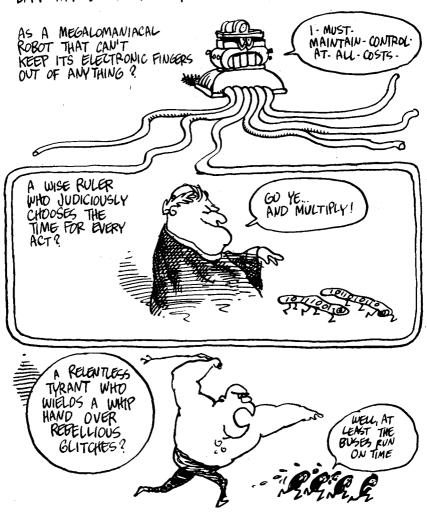
ACCORDING TO SIGNALS PASSED ALONG THE CONTROL BUS, ADDRESSES AND DATA GET ON AND OFF THE DATA ADDRESS BUS, WITH THE PROVISO THAT ONLY ONE "PASSENGER" CAN RIDE THE BUS AT A TIME



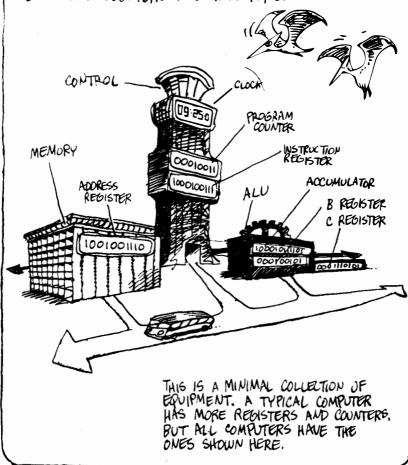


NOTE THAT ALL THE ARROWS ON THE CONTROL BUS POINT AWAY FROM THE CONTROL BECTION.

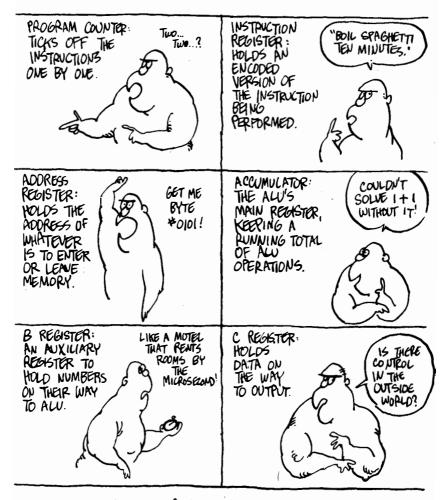
HOW ARE WE TO IMAGINE THIS CONTROL, FROM WHICH ALL DARK ARROWS POINT AWAY ??



LIKE ANYONE ELSE, CONTROL REVEALS ITS CHARACTER BY ITS BEHAVIOR ... SO LET'S FOLLOW WHAT HAPPENS IN THIS OVERSIMPLIFIED COMPUTER, WHICH FLESHES OUT THE DIAGRAM OF TWO PAGES BACK WITH SOME ESSENTIAL COUNTERS AND REGISTERS.



HERE'S WHAT THEY'RE FOR:





IN FACT, CONTROL SPENDS MOST OF ITS TIME JUST MOVING THE CONTENTS OF THESE REGISTERS AROUND!

TO SEE HOW CONTROL WORKS, LET'S FOLLOW WHAT HAPPENS WHEN THE COMPUTER ADDS TWO NUMBERS OUR VERY FIRST PROGRAM!



LIKE EVERYTHING ABOUT COMPUTERS, PROGRAMS CAN BE DESCRIBED AT VARIOUS LEVELS. WE DEGIN WITH

ASSEMBLY LANGUAGE,

WHICH SPECIFIES THE COMPUTER'S ACTUAL MOVES, BUT OMITS THE FINE DETAILS. AT THIS LEVEL, HERE'S HOW TO ADD TWO NUMBERS:

- O. LOAD THE FIRST NUMBER INTO THE ACCUMULATOR.
- 1. ADD THE SECOND NUMBER (HOLDING THE SUM IN THE ACCUMULATOR).

2. OUTPUT THE CONTENTS OF THE ACCUMULATOR.

3. HALT.



TO EXPRESS THIS IN PROPER ASSEMBLY LANGUAGE, WE MUST SPECIFY THE PRECISE LOCATION IN MEMORY OF THE TWO NUMBERS TO BE ADDED, AND CONDENSE THE WORDY STATEMENTS INTO MNEMONIC* ABBREVIATIONS. SUPPOSE, FOR EXAMPLE, THAT THE NUMBERS ARE STORED AT ADDRESSES IE AND IF (HEXADECIMAL). OUR PROGRAM BECOMES:

A TRUE ASSEMBLY: LANGUAGE PROGRAM!

O. LDA 1E

("LOAD ACCUMULATOR WITH CONTIENTS OF IE")

1. ADD IF

("ADD CONTENTS OF IF")

ST ST

2. OUT

("OUTPUT CONTENTS OF ACCUMULATOR.")

3. HALT

* MNEMONIC = MEMORY - AIDING



THE OPERATION,
WHICH DESCRIBES THE
STEP TO BE PERFORMED

THE OPERANDO WHICH GIVES THE ADDRESS ON WHICH THE OPERATOR ACTS

AS IN:
PERFORM AN
APPENDECTOMY
ON THE RESIDENT

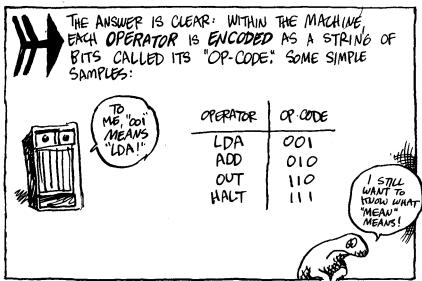
LDA 1E



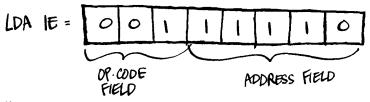
NOTE HOWEVER! SOME OPERATORS DON'T NEED AN EXPLICIT OPERAND. "OUT", FOR INSTANCE, IS UNDERSTOOD TO MPLY TO THE ACCUMULATOR.



NOW THAT WE HAVE AN ASSEMBLY-LANGUAGE PROGRAM, HOW DO WE FEED IT TO THE MACHINE—WHICH ONLY UNDERSTANDS O'S AND I'S ?



THEN A MACHINE INSTRUCTION CONSISTS OF AN OP-CODE SEGMENT, OR "FIELD," FOLLOWED BY AN ADDRESS FIELD GIVING THE OPERAND IN BINARY:



SO HERE'S OUR PROGRAM TRANSLATED INTO MACHINE LANGUAGE:

O. 1.DA 1E 001 11110 1. ADD 1F 010 11111

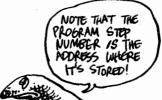
2. OUT 110 XXXXX

ANY 5 BITS ARE O.K. FOR THESE ADDRESS 3. HALT FIELDS, AS THEY'LL BE VONORED!

(ASSUMING AN INPUT DEVICE)

THE PROGRAM STEPS ARE READ INTO CONSECUTIVE MEMORY ADDRESSES, BEGINNING WITH O. THE CONTENTS OF MEMORY ARE THEN

CONTENTS
001 11110
010 11111
110 00000
111 00000

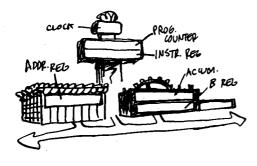


AND WE ALSO NEED TO ENTER THE DATA : THE TWO NUMBERS TO BE ADDED. ANY TWO NUMBERS WILL DO, SAY 5 AND 121. THEY GO IN ADDRESSES IE AND IF:

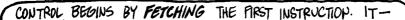
15 00000101 IF 01111001



HOW CAN THE COMPUTER DISTINGUISH DATA FROM INSTRUCTIONS? BY ASSUMING EVERYTHING IS AN INSTRUCTION, UNLESS INSTRUCTED TO DO OTHERWISE I



ONCE THE PROBRAM IS STORED, CONTROL CAN BEGIN EXECUTION, IN A SERIES OF EVEN MORE PRIMITINE STEPS CALLED MICROINSTRUCTIONS, ONE MICROINSTRUCTIONS OCCURRING WITH EACH CLOCK PULSE. ARE YOU READY FOR THE GORY DETAILS?



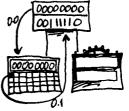
0.0. MOVES CONTENTS OF PROGRAM COUNTER (0000000 TO BEGIN WITH)

To

ADDRESS REGISTER

O.1 MOVES CONTENTS OF THAT MEMORY ADDRESS

to instruction register



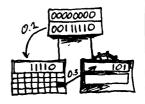
THE INSTRUCTION REGISTER NOW HOLDS THE FIRST INSTRUCTION. CONTROL "READS" IT AND—

0.2. Moves the Instruction REDISTRY'S ADDRESS FIELD

To

ADDRESS REGISTER

0.3. MOVES CONTENTS TO ACCUMULATOR OF THAT MEMORY ADDRESS



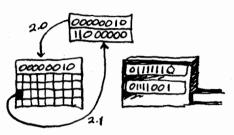
THE ACCUMULATOR IS NOW LOADED WITH THE FIRST PIECE OF DATA. ONE MICROINSTRUCTION REMAINS:

0.4 INCREMENT PROGRAM COUNTER

AND THAT'S STEP O. A BIT CONFUSED? LET'S GO THROUGH IT AGAIN WITH THE NEXT STEP, ADD. AGAIN CONTROL BEGINS WITH A "FETCH PHASE": 1.0 MOVE CONTENTS ADDRESS 00000001 OF PROGRAM To register 0101111 COUNTER (NOW 00000001) 1.1 MOVE CONTENTS TO INSTRUCTION OF THAT ADDRESS register THE INSTRUCTION IN THE INSTRUCTION REGISTER, OIO IIIII, CAUSES CONTROL TO: 1.2 MOVE ADDRESS ADDRESS To FIELD FROM 0000000 REGISTER INSTRUCTION. 01011111 REGISTER 1.3 MOVE CONTENTS B REGISTER To OF THAT MEMORY 011100 **PODRESS** 1.4 SIGNAL THE ALU TO ADD IN ACCUMULATOR 0111110 01111001 AND PUT THE AGMN, THERE'S ONE MORE STEP: good News!! It 6 1.5 INCREMENT GETS NO PROGRAM worse COUNTER

WELL, LUCKILY THE LAST TWO INSTRUCTIONS ARE EASIER:

20 AND 2.1 ARE THE SAME PETCH INSTRUCTIONS AS BEFORE, PUTTING Instruction 2 ("out") IN THE INSTRUCTION REGISTER:



THIS OP-CODE (110) CAUSES COUTROL TO -

2.2. MOVE CONTENTS OF ACCUMULATOR

TO C REGISTER



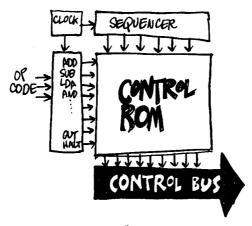
2.3. INCREMENT PROGRAM COUNTER

FINALLY, CONTROL PETCHES THE INSTRUCTION 111 ("HACT"), WHICH CAUSES CONTROL TO -

3.2 DO NOTHING



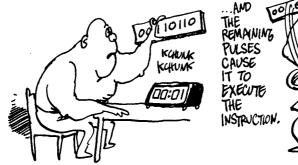
WITHOUT TOO MANY DETAILS, YOU CAN THINK OF CONTROL ROUGHLY LIKE THIS:



ITS INPUT
CONSISTS OF
CLOCK PULSES
AND OP-CODES.
ITS OUTPUT
CONSISTS OF A
SEQUENCE
OF SIGNALS TO
THE REGISTERS,
COUNTERS, ALU,
AND MEMORY.

THE "MICROPROGRAM," WHICH CONNECTS THE INPUTS TO THE PROPER OUTPUT COMBINATIONS, IS STORED IN A READ-ONLY MEMORY DEDICATED STRICTLY TO THIS PURPOSE.

THE FIRST COUPLE OF CLOCK PULSES CAUSE CONTROL TO FETCH AN INSTRUCTION...





IN REAL LIFE THE SITUATION IS MORE COMPLICATED IN DETAIL BUT THE SAME IN PRINCIPLE. THERE ARE MORE REGISTERS, AND OP-CODES ARE LONGER THAN THREE BITS, ALLOWING CONTROL TO RESPOND TO A MUCH LARGER SET OF INSTRUCTIONS. HERE'S THE INSTRUCTION SET OF A GENUINE PROCESSOR, THE MOTORPIA 6800.

BRANCH ARITHMETIC BRANCH ADD BRANCH IF ZERO ADD WITH CARRY BRANCH IF NOT ZERO SUBTRACT branch if equal SUBTRACT WITH CARPY BRANCH IF NOT EQUAL INCREMENT BRANCH IF CARRY DECREMENT BRANCH IF NO CAPPY COMPARE BRANCH IF POSITIVE PAMUL IF NEGATIVE NEGATE nch if overflow RANCH IF NO OVERFLOW LOGICAL BRANCH IF GREATER THAN branch if breater than of equal branch if less than branch if less than of equal AND OR exclusive or branch if higher branch if not higher NOT SHIFT RIGHT SHIPT LEFT BRANCH IF LOWER SHIFT RIGHT APITHMETIC BRANCH IF NOT LOWER POTATE RIGHT ROTATE LEFT SUBROUTINE CALL TEST CALL SUBROUTINE DATA TRANSFER SUBROUTINE RETURN RETURN PROM SUBPOUTINE RETURN PROM INTERPUPT LOND Store move MISCELLANGOUS NO OPERATION CLEAR POP CLEAR CARRY CLEAR OVERPLOW WAIT SET CARPY ADJUST DECIMAL ENABLE INTERRUPT SET OVERPLOW DISABLE INTERPUT BREAK

ONE GROUP OF THESE INSTRUCTIONS DESERVES SPECIAL MENTION: THE BRANCH, OR JUMP, INSTRUCTIONS.

AS WE'LL SEE,
THESE GIVE
THE COMPUTER A
LOT OF ITS
"INTELLIGENCE."
THEIR EFFECT IS TO
TRAINSTEER
CONTROL
TO ANOTHER PART OF
THE PROGRAM. THE
SIMPLEST JUMP
INSTRUCTION IS JUST
PLAIN JUMP," AS IN:



"JMP 123" CAUSES CONTROL TO ENTER 123 IN THE PROGRAM COUNTER... AND PROCEED WITH THE PROGRAM PROM THERE.

EVEN "SMARTER"

ARE CONDITIONAL
JUMPS. THEY
TRANSPER CONTROL

IF SOME
CONDITION IS
SATISFIED: FOR
INSTANCE, "JUMP
IF ZERO" MEANS
JUMP IF THE
ACCUMULATOR
HOLDS O.





IF YOU REALLY WANT TO IMAGINE THE CONTROL SECTIONS PERSONALITY. THINK OF A PERFECTLY EFFICIENT BUREAUCRAT, ACTING IN STRICT OBCHENCE TO THE COMPUTER'S REAL BOSS: THE PROBLEMAN O





IF PROGRAMS REALLY RULE THE COMPUTER, THEY DESERVE A PROPER SCIENTIFIC NAME... SOMETHING IN GREEK OR LATIN, PREFERABLY ...

> TECHNICALCULUS? REGULA RATIONOCEROUS? CEPHALONEURALGIA?



. BUT THAT'S NOT HOW IT IS IN COMPUTER SCIENCE... INSTEAD, PROGRAMS IN GENERAL ARE CALLED ROBE TO DISTINGUISH THEM FROM THE CIRCUIT BOAFDS, CATHODE

RAY MODITORS, DISK DRIVES, KEYBOARDS, AND OTHER ITEMS OF COMPUTER HARDWARB.







SOFTWARE



TUPPERWARE

WHAT'S REALLY
FUNNY ABOUT THE
NAME IS THAT
SOFTWARE IS ONE
OF THE HARDEST
THINGS ABOUT
COMPUTING!



WHILE HARDWARE HAS BEEN DROPPING IN PRICE AND GROWING IN POWER, SOFTWARE ONLY GETS MORE HORRENDOUSLY COMPLEX!





WE SEE SMALLER AND
SMALLER CHIPS WITH
BIGGER + BIGGER MANUALS!

It'S OFTEN
IMPOSSIBLE TO
ESTIMATE HOW MUCH
TIME, MONEY, AND
AGONY A GIVEN
SOFTWARE PROBLEM
WILL COST TO
SOLVE... WHAT
A WAY TO
RUN A BUSINESS!

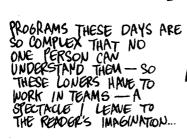


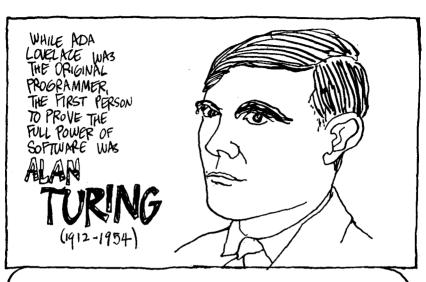
LIKEWISE THERE'S A DIFFERENCE BETWEEN THE IMAGE OF HARDWARE AND SOFTWARE WORKERS —



HARDWARE TYPES ARE ENGINEERS... INTO GADGETS... MOSTLY MEN... BOUND BY THE LAWS OF PHYSICS...

PROGRAMMERS HAVE NO TOOL BUT THEIR BRAINS...THEY'RE MORE OFTEN WOMEN... SUPPOSED TO BE SOLITARY DREAMERS WHOSE IDEAS HAVE NOTHING TO DO WITH THE LAWS OF PHYSICS!!





TURING, WHO ENJOYED LONG. DISTANCE RUNNING BACK WHEN THAT WAS CONSIDERED WEIRD, PROBABLY WENT INTO COMPUTERS TO SHRINK THE SIZE OF HIS JOGGING CLOCK.



TURING MACHINES

APEN'T REPL MACHINES...

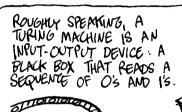
THBY'RE ABSTRACT

MACHINES, EXISTING ONLY

IN THEORY...



SOFTUARE ENGLIEGE'S DREAM — NO HARDWARE!



THE OUTPUT CONTINE
DEPENDS ONLY ON THE
PRESENT INPUT (O OR 1)
AND THE PREVIOUS OUTPUT.

THE NATURE OF THE OUTPUT IS UNIMPORTANT.

THE MAIN THING IS
THAT THE CHANGES
FROM ONE OUTPUT
STATE TO THE NEXT
ARE GIVEN BY
DEPINITE RULES,
CALLED THE
TRANSITION
RULES.

THERE'S A

TURING MACHINE

THAT CAN

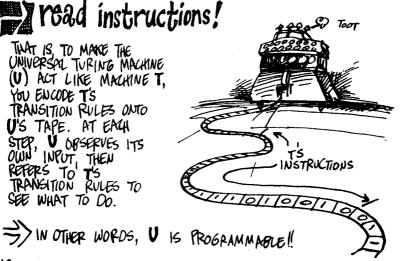
ADD!

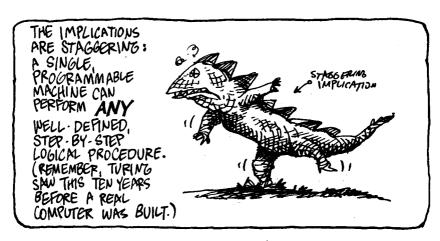
THE REASON TUPHING
MACHINES ARE IMPORTANT
IS THAT THEY ARE A
A WAY OF THINKING
PHYSICALLY ABOUT LOCK.
ANY WELL DEFINED, STEP-BY-STEP
LOCKAL PROCEDURE
CAN BE EMBODIED
IN SOME TURING MACHINE.

+FOR DETAILS, SEE J. WEIZENBRUM'S COMPUTER POWER MO NUMBER REASON, CAMPTER 2.

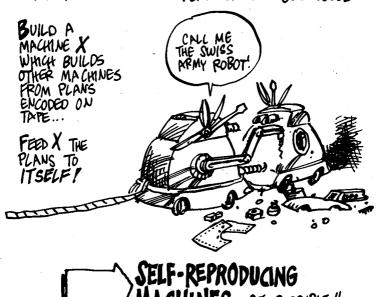


THE TRICK IS THAT THE UNIVERSAL TURING MACHINE CAU...



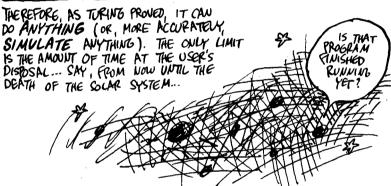


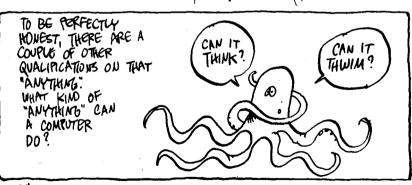
JOHN VON NEUMANN CARRIED TURING'S IDEMS A STEP FURTHER VON NEUMANN REALIZED THAT ONE COULD:







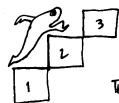




IN A WORD, COMPUTERS DO



AN ALGORITHM
IS SIMPLY
ANY WELL DEPINED,
STER-BY-STER
PROCEDURE: A
RECIPE, IF YOU
WILL!



STEP-BY-STEP,
MEANING EACH STEP
IS COMPILTED BEFORE
THE NEXT IS BEGUN.

WELL DEFINED,
MEANING EACH STEP
IS COMPLETELY
DETERMINED BY
CURRENT INPUT AND
THE RESULTS OF
PREVIOUS STEPS.
NO AMBIOUITY ALLOWED!



EXAMPLES OF ALGORITHMS:

"IF NUCLEAR WARHEADS ARE FALLING LINE HAILSTONES, I WILL LIE DOWN AND TRY TO ENJOY IT.

OTHERWISE, I WILL GO TO WORK AS USUAL."



LY IT'S AN ALGORITHM BECAUSE I ALWAYS KNOW WHAT TO DO:

- 1. CHECK TO SEE IF WARHEADS PRE FALLING
- 2. IF YES, LIE DOWN + ENJOY!
- 3. IF NO, GO TO WORK.



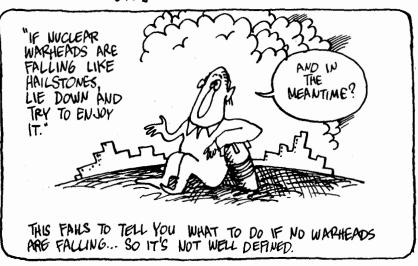
LIKEWISE, ALGERAIC FORMULAS REPRESENT ALGORITHMS $y = x^2 + 2x + 10$ MEANS—

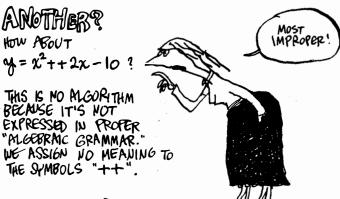
> IF YOU UNDERSTAND, LIE DOWN AND ENJOY YNTRELF!

- (1) INPUT A NUMBER &
- (2) MULTIPLY & TIMES ITSELF
- (3) MULTIPLY & TIMES 2
- (4) ADD THE RESULTS OF (1) AND (3)
 - (5) ADD 10 TO THE RESULT OF (4)



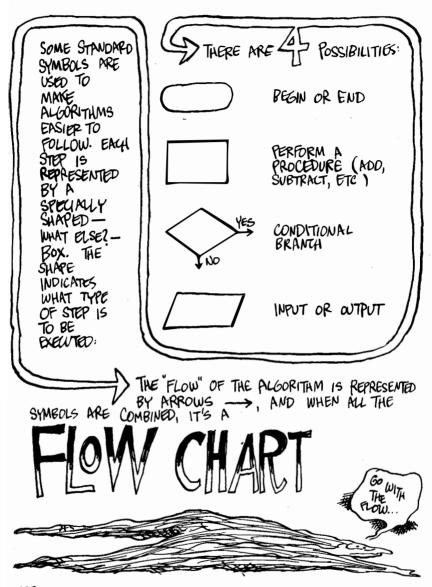
EXAMPLES OF NONDALGORITHMS:



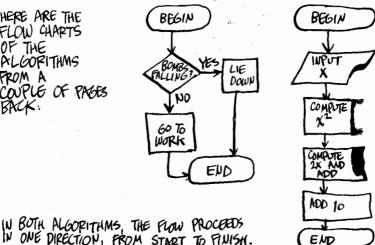




IF YOU TRY TO MAKE A COMPUTER DO A NON-ALGORITHM, IT WILL JUST SIT THERE PLASHING ERROR MESSAGES!







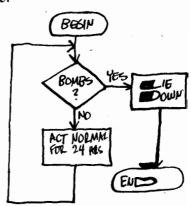
IN ONE DIRECTION, FROM START TO FINISH. IT'S ALSO POSSIBLE FOR THE FLOW OF ALGORITAMS TO JUMP FORWARD

OR BACKWARD. FOR EXAMPLE, LET'S REWRITE THAT FIRST

ALGORITUM:

1. IF BOMBS ARE FALLING. 60 To STEP 2. OTHERWISE, 60 To STEP 4.

- 2. LIE DOWN AND ENJOY!
- 3. GO TO STEP 6.
- 4. LEAD A NORMAL LIFE FOR 24 HOURS
- 5. GO TO STEP 1
- 6. END



YOU MAY FIND THE FLOW CHANDT EASIER TO GRASP THAN THE WRITTEN "PROGRAM." NOTE THAT IT MAY CONTINUE INDEPNITATELY! FLOW CHARTS ARE USEFUL
IN HELPING TO DESIGN
ALGORITHMS—SIMPLE ONES,
ANY WAY— AND DESIGNING
ALGORITHMS IS WHAT
COMPUTER PROGRAMMING
IS ALL ABOUT!!





200

LET'S TRY A COUPLE MORE EXAMPLES... A LITTLE MORE LIKE WHAT A COMPUTER MIGHT ACTUALLY BE ASKED TO DO...



MUSTIFIE PRUG-INST

THIS ONE AGKS THE COMPUTER TO EVALUATE THE EXPRESSION

X²+2×+10

NOT JUST AT OME VALUE OF X, BUT FOR MANY VALUES, NAMBLY

X=0,0.1,0.2,0.3, ... AND SO ON...

UP TO 2.0.



FOR ROOMMATE RECEIPTS" WE REASON LIKE SO:

LET S=SOPHIE'S EXPENSES L=LISA'S EXPENSES

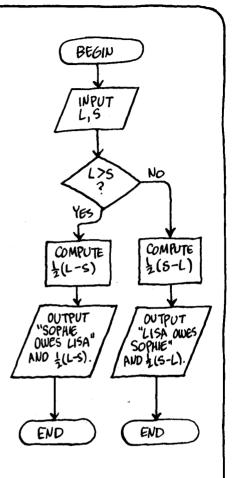
THEN THE TOTAL EXPENSE IS S+L, AND EACH RODMMATE'S SHARE IS $\frac{1}{2}(S+L)$.

IF LISA OUTSPENT SOPHIE, SO L>S*, THEN SOPHIE OWES LISA $\frac{1}{2}(S+L)-S$, OR $\frac{1}{2}(L-S)$.

OTHERWISE (WHEN SZL*), LISA OWES SOPHIE

$$\frac{1}{2}(S-L)$$
.

THE ALBORITHM'S OUTPUT IS TO TELL US WHO OWES WHOM AND HOW MUCH.



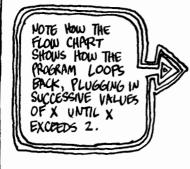
* > MEANS "IS GREATER THAN"; \geq MEANS "IS GREATER THAN OR EQUAL TO"; \leq MEANS "IS LESS THAN"; \leq MEANS "IS LESS THAN OR EQUAL TO".

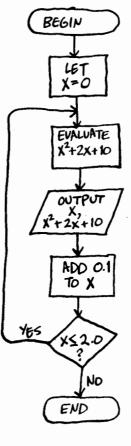
IN "MULTIPLE PLUG-(NS," WE WANT TO EVALUATE A STAGGE EXPRESSION, $x^2+2x+10$, REPEATEDLY AT DIFFERENT VALUES OF X (NAMELY 00, 0.1, 0.2, ..., 1.9, 2.0)

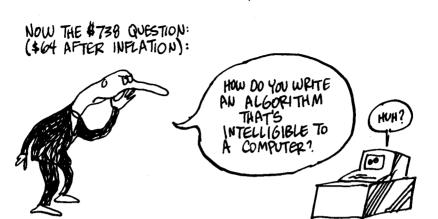
THE CORE OF THE ALGORITHM WILL BE THIS LOOP:

- 1. PLUG THE CARRENT VALUE OF X INTO X2+2x+10
- 2. PAINT THE RESULT
- 3. NEXT X
- 4. RETURN TO STEP 1.

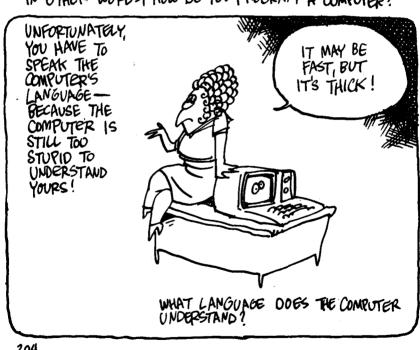
WE ALSO HAVE TO SPELLLY WHAT X TO STAFT WITH, WHEN TO STOP, AND HOW TO COMPUTE "NEXT X."







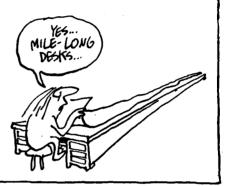
IN OTHER WORDS, HOW DO YOU PROGRAM A COMPUTER?



AT THE VERY BEGINNING, PROGRAMMERS WROTE DIFECTLY IN "MACHINE LANGUAGE" — BINARY CODE. THIS WAS OBVIOUSLY A HEADACHE!



SOON TACY
SWITCHED TO
ASSEMBLY LANGUAGE
(SEE P. 174), AIDED
BY AUTOMATIC
"ASSEMBLERS"
WHICH TRANSLATED
ASSEMBLY LANGUAGE
MNEMONICS INTO
MACHINE CODE.
STILL SOMETHING MORE
WAS NEEDED!



AND FINALLY,

SOURCE CODE

COMPILER OF
INTERPRETER

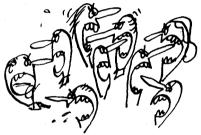
OBJECT CODE

THE HIGHER-LEVEL PROGRAMMING LANGUAGES

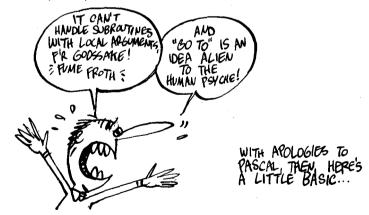
WERE INVENTED. THESE CONTAIN FAMILIAR ENGLISH-LIKE COMMANDS, SUCH AS "PAINT," "READ," AND "DO," WHICH ARE TRANSLATED INTO MACHINE LANGUAGE BY COMPLEX PROGRAMS CALLED COMPILERS OR INTERPRETEDS. HIGHER-LEVEL PROGRAMS ARE SOMETIMES CALLED "SOURCE CODE," AND THE MACHINE-LANGUAGE TRANSLATION IS CALLED "OBJECT CODE."

THE FIRST MIGHER LEVEL
LANGUAGE WAS
FORTRAN ("FORMULA
TRANSLATOR"), WHICH MADE
ITS DEBUT IN THE EARLY
1950'S. SINCE THEN,
LITERALLY HUNDREDS OF
LANGUAGES HAVE BEEN
WRITTEN, ENGH WITH ITS
OWN ARMY OF RABID
DEVOTEES!





WE'RE GOING TO TAKE A QUICK LOOK AT BASIC—
BEGINNER'S JUL-PURPOSE SYMBOLIC [INSTRUCTION GODE.
BASIC IS EASY TO LEARN AND WIDELY USED, DESPITE
CRITICISM (ESPECIALLY BY PASCAL ADMIRERS) THAT IT PROMOTES
"BAD PROGRAMMING HABITS."







THERE ARE TWO
WAYS TO WRITE A
BASIC PROGRAM:
WITH PENCIL AND
PAPER, OR DIRECTLY
AT THE COMPUTER.

It's GOOD PRACTICE TO PLAN PROGRAMS ON PAPER FIRST, TO WORK OUT THE ESSENTIAL IDEAS AND STRUCTURE, BUT EVENTUALLY YOU MUST SIT DOWN AT THAT KEYEDARD!



Some Machines afe ready for basic as soon as you turn them on. others only bring it up on command. If In Doubt, ask!



WHEN THE COMPUTER IS READY IT GIVES YOU A "PROMPT" OF SOME KIND: THE WORD "READY" OR JUST THE SIGN ">".



THE COMPUTER KEYBOARD RESEMBLES A STANDARD TYPE-WRITER'S "QWERTY" KEYBOARD ... EXCEPT THAT AS YOU TYPE, CHARLACTERS APPEAR ON THE CRT (CATHODE RAY TUBE) SCREEN, INSTEAD OF ON PAPER. TO GO TO THE NEXT LINE, HIT THE RETURN (2) KEY. HERE'S A SIMPLE BASIC PROGRAM:

10 REM BASIC MULTIPLICATION

20 READ A, B

30 DATA 5.6, 1.1

40 LET C=A*B

50 PRINT "THE PRODUCT IS"; C

60 END

THE PROGRAM IS NOW STORED IN MEMORY. TO RUN IT, TYPE "RUN", FOLLOWED BY THE RETURN KEY.

THE SCREEN DISPLAYS:

RUN THE PRODUCT IS 6.16

BASIC MATH: A-B } AS USUAL A* B... A TIMES B A/B ... A DIVIDED BY B ATB... A TO THE BM POWER



POW POINTS TO NOTE:



- DEVERY LINE BEGIND WITH A LINE NUMBER (10, 20,...).
 EVERY LINE OF A BASIC PROGRAM MUST HAVE A NUMBER! IT'S WISE TO COUNT BY TENS, SO YOU CAN INSERT LINES LATER.
- THE FIRST LINE (10) IS A REMARK. REMARKS EXPLAIN THE PROGRAM BUT AREN'T EXECUTED BY THE COMPUTER. THE PREFIX "REM" IDENTIFIES REMARKS. WE MIGHT INSERT ONE HERE:

20 READ A, B 25 REM THESE ARE THE #5 TO BE MULTO 30 DATA 5.6, 1.1

PROGRAM STATEMENTS CONSIST OF INSTRUCTIONS ("LET", ETC), NUMBERS (5.6, 1.1), VARIABLES (A,B,C), TEXT ("THE PRODUCT IS"), AND PUNCTUATION.

50 PRINT "THE PRODUCT IS"; C

SEACH OF THESE HAS A PREASE







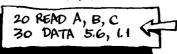
A NUMERICAL VARIABLE IN BASIC 16 LIKE A VARIABLE IN ALGEBRA IT ASSUMES A NUMERICAL VALUE, WHICH MAY VARY (BUT IT HAS ONLY ONE VALUE AT A TIME!). ONLY THESE SYMBOLS CAN BE USED AS VARIABLES:

A, B, C, D, \dots	Z
AO, BO, AND EVERYTHIN	Zo
A1, BI, EVERYTHI	^{N6} Z1
Between	ni i
A9, B9,	Z9

THERE ARE SEVERAL WAYS TO ASSIGN A VALUE TO A VARIABLE:

20 READ A, B
30 DATA 5.6, 1.1
COMMAS ARE BESENTIAL!

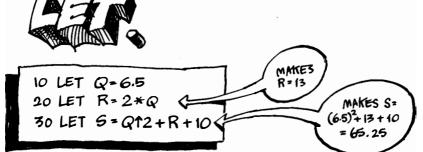
THIS INSTRUCTS THE COMPUTER TO ASSIGN THE NUMERICAL VALUES IN THE DATE: STATEMENT - IN ORDER - TO THE VARIABLES IN THE REPORT STATEMENT.



THIS IS A BUG!



ANOTHER WAY TO ASSIGN VALUES TO VARIABLES IS WITH



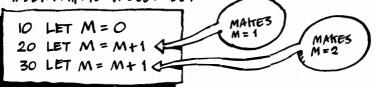
THE LET STATEMENT ASSIGNS THE VALUE ON THE RIGHT OF THE EQUALITY SIGN, "=", TO THE VARIABLE ON THE LEFT. THE RIGHT-HAND SIDE MAY BE A NUMBER, OR SOME MATHEMATICAL EXPRESSION INVOLVING OTHER VARIABLES— HE LONG AS THEY PAREADY HAVE VALUES!!

10 LET Q=6.5

20 LET Q=0.5*R

30 LET 5= Q12+R+10

HERE STATEMENT 20 DOES NOT ASSIGN ANY VALUE TO R, BECAUSE R IS NOT ON THE LEFT SIDE OF "=". IN FACT, IF R HASN'T BEEN ASSIGNED SOME VALUE EARLIER IN THE PROGRAM, THEN STATEMENT 20 GIVES Q AN INDETERMINATE VALUE! BUT—



THESE STRANGE-LOOKING STATEMENTS ARE PEFECTLY O.K! "LET M = M + 1" MEANS "ASSIGN TO THE VACHABLE M A VALUE EQUAL TO ITS CURRENT VALUE PLUS 1."



THIS IS AN OUTPUT COMMAND, MEANING "DISPLAY ON THE SCREEN," NOT "PRINT ON PAPER."



YOU CAN PRINT ANY TEXT:

10 PRINT "ANY NUKES TODAY?"

ANY NUKES TODAY?

QUOTATION MARKS ESSENTIAL!

QUOTATION MARKS REMOVED

PRINT A VARIABLE AND YOU GET ITS VALUE:

10 LET X=77001 20 PRINT X

FUN

77001

BUT-

10 LET X= 77001 20 PRINT "X" (QUOTATION MARKS MAKE THE COMPUTER TREAT X AS A TEXT.

PRINT A MATHEMATICAL EXPRESSION AND YOU GET ITS VALUE:

10 LET Z=1.5 20 PRINT Z+2+2*Z+10

RUN 15.25

X

BECAUSE

 $(1.5)^2 + 2 \times 1.5 + 10$

SEMICE CSD

A SEMICOLON AFTER A PRINT STATEMENT CAUSES THE NEXT PRINT STATEMENT TO DISPLAY ITS OUTPUT ON THE SAME LINE AND DIRECTLY AFTER THE FIRST ONE'S:

10 LET A = 1
20 PRINT "INFINITY IS MORE THAN";
30 PRINT A
RUN
INFINITY IS MORE THAN 1

IT'S O.K. TO ABBREVIATE THIS:

10 LET A=1
20 PRINT "INFINITY IS MORE THAN"; A
RUN
INFINITY IS MORE THAN 1

FOR EXAMPLE, WE COULD REWRITE THE PROGRAM ON P. 208.

IO REM BASIC MULTIPLICATION

10 READ A, B

30 DATA 5.6, 1.1

40 LET C= A*B

50 PRINT "THE PRODUCT OF"; A; "AND"; B; "Is"; C;"."

60 END

PUN

THE PRODUCT OF 5.6 AND 1.1 IS 6.16.

THERE ARE ALSO SOME HIFTY TRICKS USING THE GOMMENT AND PRINT, BUT WE WON'T GET INTO IT...



THE STATEMENT ALLOWS THE USER TO ASSIGN VALUES TO VARIABLES WHILE THE PROGRAM IS RUNNING.



THE FORM OF THE STATEMENT:

INPUT A

WHEN THE PROGRAM RUNS AND REACHES AN INPUT STATEMENT, THE SCREEN DISPLAYS:

?

THIS INDICATES THAT THE PROGRAM HAS HALTED, AWANTING INPUT. YOU TYPE SOME NUMBER (FOLLOWED BY "RETURN," AS ALWAYS!):

5.6

AND THE PROGRAM CONTINUES PUNISHED.
"INPUT" AND "PRINT" CAN BE USED IN COMBINATION TO LET YOU KNOW WHAT SORT OF INPUT IS EXPECTED:

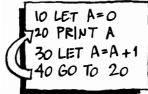
10 BASIC DIVISION
20 PRINT "TYPE THE NUMERATOR."
30 INPUT N
40 PRINT "TYPE THE NON-ZERO DENOMINATOR."
50 INPUT D
60 PRINT N; "/"; D; "="; N/D
70 END
RUN
TYPE THE NUMERATOR.
? 5
TYPE THE NON-ZERO DENOMINATOR.
? 8
5/8 = 0.625

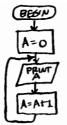
TYPED BY THE USER.

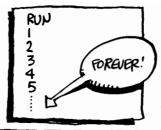


TAIS IS THE UNCONDITIONAL BRANCHING INSTRUCTION.

"GO TO (LINE NUMBER)" TRANSFERS CONTROL TO A LINE OTHER THAN THE NEXT. THE PROGRAM THEN CONTINUES FROM THERE, AS IN THIS ENDLESS LOOP:



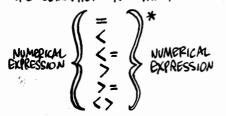






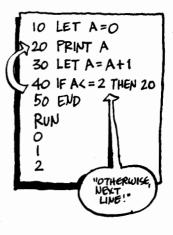
IS THE "SMART," CONDITIONAL JUMP.

IT HAS THE GENERAL PORM IF (CONDITION) THEN (LINE NUMBER). THE CONDITION HAS THE FORM:



AS IN IF A C=B THEN 30

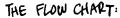
THIS ALWAYS INCLUDES THE UNSTATED INSTRUCTION, "OTHERWISE, GO TO THE NEXT LINE."

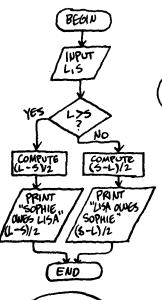


^{* &}lt; LESS THAN, <= LESS THAN OF EQUAL TO, > GREATER THAN OF EQUAL TO; <> DOES NOT EQUAL.

THIS IS ENOUGH TO WRITE BASIC PROGRAMS FOR THE TWO ALGORITHMS FROM P. 201:







THE PROGRAM:

IO PRINT "LISA SPENT"

20 INPUT L

30 PRINT "SOPHIE SPENT"

40 INPUT S

50 IF L."> 5 THEN 80

60 PRINT "LISA ONES SOPHIE"; (S-L)/2

70 GO TO 90

80 PRINT "SOPHIE OWES LISA"; (L-S)/2

7190 END

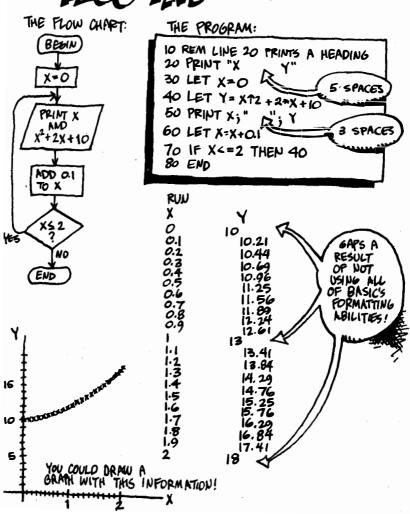
SEE NOW "IF-THEN" AMD "60 TO" ARE USED? IF L>S, THEN LINES 60 AMD TO ARE NOT EXECUTED.
OTHERWISE, THEY ARE EXECUTED, AND LINE TO ENSURES THAT LINE 80 IS SKIPPED.

IF THE PROGRAM IS RUN:

RUN LISA SPENT ? 93.75 SOPHIE SPENT ? 77.38 SOPHIE OWES LISA 8.185

NOW WE



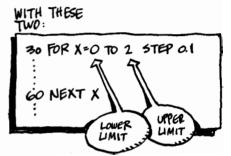


THE "MULTIPLE PLUG-INS"
LOOP IS SO TYPICAL
THAT ALL PROGRAMMING
LANGUAGES HAVE
SPECIAL COMMANDS JUST
FOR SUCH REPETITIONS.
IN BASIC, IT'S

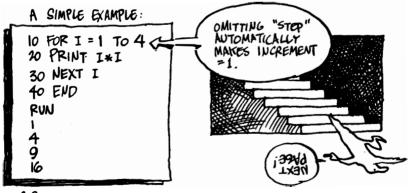




30 LET X=0 :: 60 LET X=X+0.1 70 IF X<=2 THEN 30



THE STATEMENT INITIALLY SETS THE VARIABLE EQUAL TO THE LOWER LIMIT, EXECUTES THE LIMES UP TO "NEXT," INCREMENTS THE VARIABLE BY THE AMOUNT "STEP," AND REPEATS THE LOOP UNTIL THE UPPER LIMIT IS EXCEDED.



HE BANG



PROBLEMS? WHO HAS PROBLEMS?

1. WHAT DOES THIS PROGRAM 00?

IO INPUT N
20 FOR I=1 TO N
30 PRINT I*I
40 NEXT I
50 END

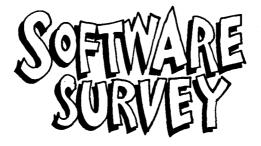
- 2. REWRITE THE "MULTIPLE PLUG-INS" PROGRAM USING THE "FOR NEXT" STATEMENT.
- 3. WRITE A PROGRAM WHICH ADDS THE INTEGERS (WHOLE NUMBERS) FROM 1 TO 1,000,000.
 DITTO FROM 1 TO N, FOR AMY N.
- 4. IN THE FIBONACCI SEQUENCE 0, 1, 1, 2, 3, 5, 8, 13, 21, 34,...
 ENCH NUMBER IS THE SUM OF THE PREVIOUS TWO NUMBERS.
 WRITE A PROGRAM WHICH GENERATES THIS SEQUENCE.
- 5. READ ENOUGH OF A BASIC TEXTBOOK TO WRITE A "ROOMMATE RECEIPTS" PROGRAM FOR ANY NUMBER OF ROOMMATES.

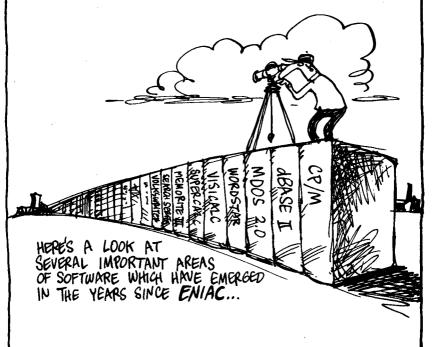


THERE ARE PLENTY OF OTHER BASIC FEATURES, ENOUGH TO FILL ENTIRE BOOKS — AND IN FACT TONS OF BOOKS ON BASIC HAVE BEEN PUBLISHED.



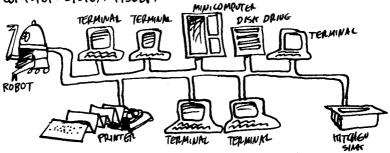




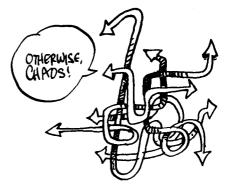




APPLICATIONS SOFTWARE DOES "REAL WORLD" JOBS, WHILE SYSTEMS SOFTWARE EXISTS PURELY TO REGULATE THE COMPUTER SYSTEM ITSELF.



A SYSTEM TYPICALLY
CONSISTS OF ONE OR
MORE INPUT/OUTPUT
DEVICES (TERMINALS,
PPLINTERS, CAPD READERS,
COMMUNICATIONS PORTS),
PROCESSORS, MEMORY
UNITS (MAIN AND MASS),
AND WHO KNAWS
WHAT ELSE.
SOMETHING HAS TO
COORDINATE IT ALL!



THE PROGRAM
THAT DOES IT
IS CALLED THE

OPERATING SYSTEM.

IF YOU THINK OF THE COMPUTER'S CORE AS A GIANT ELECTRONIC FILING CABINET (WITH A CALCULATOR ATTACHEO), THEN THE OPERATING SYSTEM

ST CREATES THE STRUCTURE

MANAGES MEMORY
SO THAT DIFFERENT
FILES DON'T BUMP
INTO EACH OTHER
AT REBULATES ACCESS

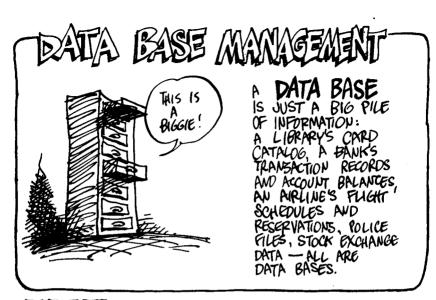
TO THE FILES AND
THE MOVEMENT OF
INFOLMATION TO AND FROM
OTHER PARTS OF THE
SYSTEM...

etc,

BESIDES THE CRENATIVE
SYSTEM, SYSTEM SOFTWARE
INCLUDES OTHER PROGRAMS
"IN THE SYSTEM," SUCH AS
LONDERS (WHICH LOAD PROGRAMS
INTO MEMORY) AND CONFILERS
(WHICH TRANSLATE HIGHER-LEVEL
LANGUNGE INTO MACKINE CODE).

ALL INVISIBLE TO THE USER!





A DATA BASE MANAGEMENT PROGRAM ORGANIZES, UPDATES, AND PROVIDES ACCESS TO THE DATA BASE.

IN THE CASE OF AN AIRLINE, FOR EXAMPLE, THE COMPUTER HAS TO BOOK RESERVATIONS, ASSIGN SEATS, ERASE RESERVATIONS WHEN THE CUSTOMER CANCELS, MAKE REASSIGNMENTS IF A PUGHT IS CANCELED, PRINT THE TICKETS, AND PROVIDE ALL THE FLIGHT INPORMATION TO TRAVEL AGENTS—WOLLD WIDE!



WORD

PROGESSUNG

A "PERSONAL" USE FOR COMPUTERS...

WORD PROCESSING
SOFTWARE ALLOWS YOU
TO WRITE, EDIT, AND
PORMAT TEXT—ALL
PROM THE SAME
KUBOARD. YOU CAN
GO FROM PIPST TO
PINAL DEAPT ELECTRONIUM,
BEFORE EVER
PRINTING A WORD.



THERE ARE ALSO
PROGRAMS TO
CORRECT SPELLINGAND EVEN TO
FIX SYNTAX AND
GRAMMAR. SOON
ILLITERATES WILL
BE CREATING
MASTER PIECES!



A SMALL COMPUTER WITH WORD PROCESSING CAN BE QUITE INEXPENSIVE... THE CATCH IS THAT A "LETTER QUALITY" PRINTER CAN COST TEN TIMES THE PRICE OF A TYPEWRITER!







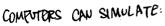
CRAY-1 COMPUTER, CAPABLE OF 100 MILLION OPERATIONS VER SECOND!

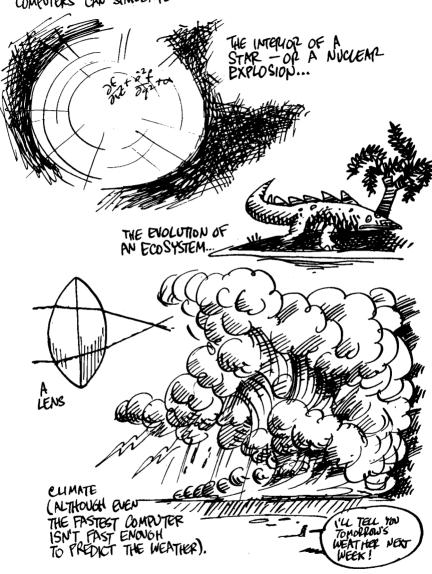
SCIENCE DEPENDS ON MATHEMATICS, AND COMPUTERS ARE SUPER MATH MACHINES.
THE FASTEST, MOST POWERFUL COMPUTERS ARE MAINLY APPLIED TO SCIENTIFIC PROBLEMS.

THESE "SUPERCOMPUTERS" EXCEL AT STIMULATION OF THE IDEA BEHIND SIMULATION IS TO FEED THE COMPUTER THE EQUATIONS GOVERNING A PHYSICAL SYSTEM AND THEN MATHEMATICALLY "MONE" THE SYSTEM ACCORDING TO THOSE LAWS.



TMIE SPACE TRAVEL: A
COMPUTER CAN GUIDE A CPAFT
TO THE MOON, BECAUSE IT
CAN INTERNALLY SIMULATE
THE ENTIRE FLIGHT!!







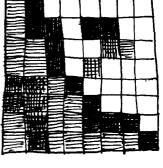


FROM THE SIMPLEST PONG" SCREEN TO THE MOST SOPHISTICATED FLIGHT SIMULATOR, THE IDEA IS THE SAME:

DIVIDE THE SQREEN AREA INTO A LARGE NUMBER OF TINY RECTANGUES ("PIXELS") AND ASSIGN EACH ONE A COLOR AND BRIGHTNESS.



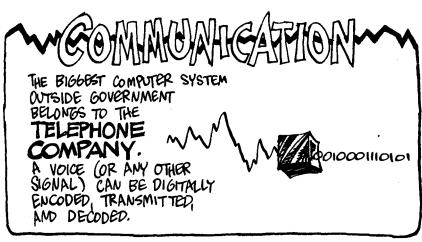
CORNERS!



BUT THERE ARE ALSO ALGORITHMS FOR SMOOTHING CORNERS!

UNFORTUNATBY, IT
TAKES A LOT OF
COMPUTER POWER TO
DO FANCY ORAPHICS.
SMALL COMPUTERS
MOSTLY DO THINGS
LIKE MAKE PIE
CHAPTS...





COMPUTERS ALSO CONTROL THE ROUTING AND SWITCHING OF
CIALLS THROUGH THE NETWORK—
AND KEEP TRACK OF
EVERYONE'S BILL!

THERE'S

THERE'S

THERE'S





DESPITE THEIR INCREDIBLE SPEED AND ACCURACY, COMPUTERS ARE LOUSY AT PATTERN RECOGNITION, ANALYSIS, HUNCH-PLAYING, AND UNDERSTANDING HUMAN LANGUAGE!

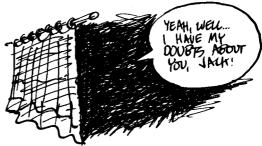


SO A BETTER QUESTION IS: HOW CAN YOU TELL IF A MACHINE IS THINKING?

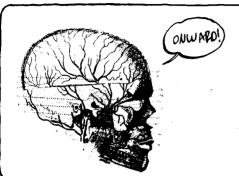
ALAN TURING SUGGESTED THIS TEST. SUPPOSE YOU COULD COMMUNICATE WITH SOMETHING, OR SOMEONE, CONCENED PROM VIEW. IF, ON THE BASIS OF THE CONVERSATION, YOU COULDN'T SAY WHETHER IT WAS MACHINE OR HUMAN, YOU WOULD HAVE TO SAY IT WAS THINKING!







I PERSONALLY DISLIKE THIS CRITERION, ON THE GROUNDS THAT A SIMULATION ISN'T THE REAL THING...



THIS PHILOSOPHICAL MUDDLE HASN'T STOPPED PEOPLE FROM TRYING TO MAKE MACHINES THINK.
THEY'VE HAD SOME SUCCESS WITH SO CALUED EXPERT SYSTEMS, WHICH MIMIC HUMAN EXPERTS IN VARIOUS FIELDS.

HOW DO YOU CHEATE AN EXPERT SYSTEM?

FIRST, INTERVIEW
A BUNGL OF
EXPERTS —
GEOLOGISTS, FOR
EXAMPLE — AND
FORCE THEM TO
SPELL OUT THE
ALLOCRITHMS BEGLIND
THEIR SKILLS, HUNTHES
AND BRAINSTORMS.



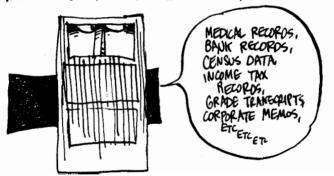
THEN LOAD THE COMPUTER'S MEMORY WITH THE HUMANS'
KNOWLEDGE BASE...
AND THE RESULT
IS (SOMETIMES)
A PROBRAM WHICH
CAN OUTPERFORM
ANY HUMAN!!

CRYPTOGRAPHY GHH!

THERE ARE STANDARD
ODES LIKE ASCII
(P. 128) FOR
ODIVERTING WRITTEN
TEXT INTO DINARY...
BUT WHAT AROUT USING
COMPUTERS FOR
SECRET
CODES ??



SECRET CODES USED TO BE STRICTLY MILITARY AND SPY STUFF, BUT HOW MORE AND MORE SENSITIVE INFORMATION IS STORED IN COMPUTER SYSTEMS:





SCRAMBLING DATA HAS BECOME AN IMPORTANT WAY OF PROTECTING PRIVACY !!

ORDINARILY, INFORMATION IS STORED AS A BINARY STRING ANY COMPUTER CAN READ: THE **PLAINTEXT**, IN CRYPTOGRAPHIC JARGON. TO ENCRYPT IT YOU APPLY SOME ALGORITHM **S**, WHICH CONVERTS IT TO A SCRAMBLED MESSAGE CALLED THE **CYPHERTE**XT.

PLAINTEXT S CYPHERTEXT

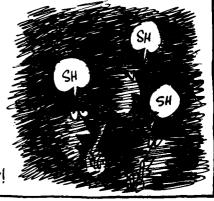
THEORETICALLY, IT'S IMPOSSIBLE TO RECONSTRUCT THE PLAINTEXT FROM THE CYPHERTEXT WITHOUT KNOWING SOMETHING ABOUT S ... HOWEVER, A POTENTIAL CODE-BREAKER COULD PUT A COMPUTER TO WORK SEARCHING FOR S.

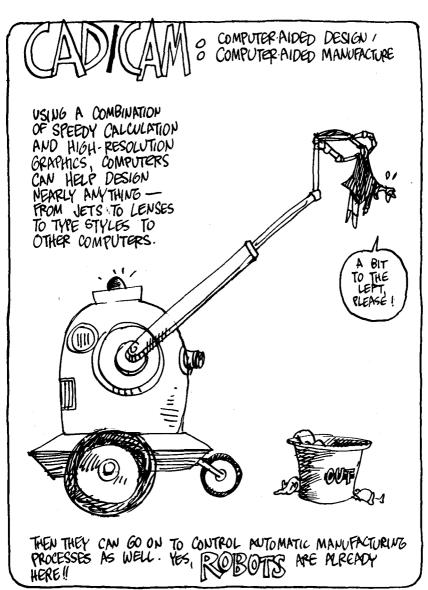


TO BE SECURE, S HAS TO BE SO COMPLICATED THAT EVEN THE PASTEST COMPUTER WOULD TAKE, SAY, A FEW MILLION YEARS TO FIGURE IT OUT!



RECENTLY, THE NATIONAL BUREAU OF STANDARDS
APPROVED A FAMILY OF ALGORITHMS AS A DATA ENCRYPTION STANDARD FOR THE NATION.
SEVERAL SCIENTISTS SUSPECT THAT THIS STANDARD IS JUST COMPLEX ENOUGH TO STYMIE ORDINARY COMPUTERS, BUT NOT TOO TOUCH FOR THE NINE ACRES OF COMPUTERS OF THE NATIONAL SECURITY AGENCY!







THE MILITARY CAN USE JUST ABOUT EVERY TYPE OF SOFTWARE WE'VE MENTIONED — AND THEN SOME!



FLIGHT SIMULATORS CAN TRAIN PILOTS RIGHT ON THE GROUND...



GREAT GRAPHICS ON THESE!

THEN THERE ARE THE WHICH CAN FOLLOW A MISSILES,

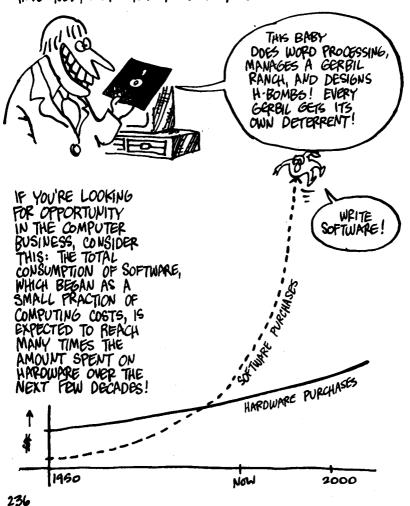


1'M ABOUT TO BLOW MYSELF SMART IS THAT?

... NOT TO MENTION DATA
PROCESSING AND CRYPTOGRAPHY...
SO CREAT IS THE DEFENSE
DEPARTMENT'S SOPTWARE
NEED THAT THEY HAVE
THEIR OWN PROGRAMMING
LANTOVAGE: ADA, NAMED
APTER THE UNFORTUNATE
LADY LOVELACE.



THIS LITTLE SURVEY ONLY BEGINS TO SUGGEST THE RANGE OF SOFTWARE CURRENTLY AVAILABLE. EVERY DAY THERE'S MORE... SOME PROGRAMS MOVE INTO NEW AREAS, WHILE OTHERS INTEGRATE EXISTING ROUTINES INTO NEW, MORE POWERFUL PACKAGES.



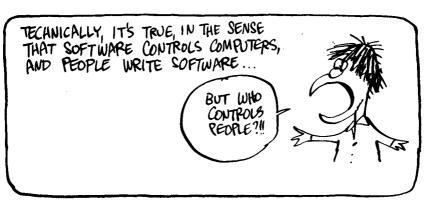
in conclusion,

A FEW WORDS ABOUT THIS FAMILIAR SENTENCE:

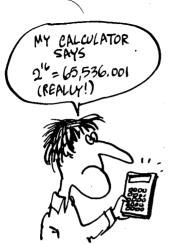
COMPUTERS
ONLY DO WHAT
PEOPLE TELL
THEM TO DO!



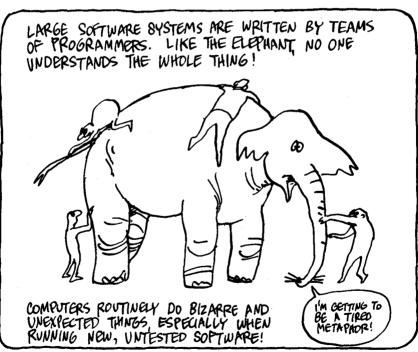
(which is what computer scientists Say when they want to be reassuring...)



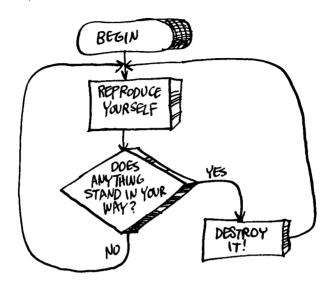




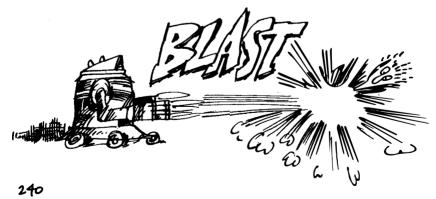
ANOTHER PROBLEM IS THAT ALGORITHMS DON'T ALWAYS DO EXACTLY WHAT THEY ARE SUPPOSED TO:



FINALLY, CONSIDER THIS OMINOUS ALGORITHM:



WHILE NO COMPUTER IS INTELLIGENT, MOBILE, OR WELL EQUIPPED ENOUGH — YET — TO EXECUTE THESE INSTRUCTIONS, SUCH A MACHINE REMAINS A THEORETICAL POSSIBILITY.
THIS PROGRAM WOULD MAKE IT SOMETHING VERY MUCH LIKE A COMPETING LIFE FORM!!!

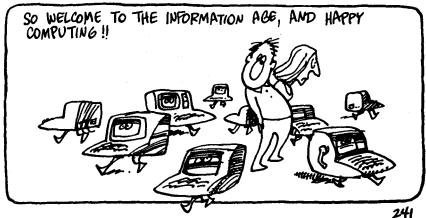


AND IF YOU THINK THAT
BECAUSE "IT'S ONLY A
MACHINE," YOU CAN ALWAYS
TURN IT OFF, PONDER
THE WORDS OF NORBERT
WIENER, A SCIENTIST WHO
THOUGHT DEEPLY ABOUT
THESE THINGS:



"TO TURN A MACHINE OFF EFFECTIVELY, WE MUST BE IN possession of INFORMATION AS TO WHETHER THE DANGER POINT HAS COME. THE MERE FACT THAT WE HAVE MADE THE MACHINE DOES NOT GUARANTEE THAT WE shall have the proper Information to 00 THIS....THE VERY SPEED OF... MODERN DIGITAL machines stands in THE WAY OF OUR ABILITY TO PERCEIVE AND THINK THROUGH THE THE INDICATIONS OF DANGER." *

> * CYBGRUETICS, SECOND EDITION, 8-175



SOME FURTHER READING:



MEDIEVAL AND EARLY MODERN SCIENCE BY A.C. CROMBIE. TELLS HOW

ISLAMIC SCIENCE CAME TO EUROPE.

THE MAKING OF THE MICRO BY C. EVAUS. NICE DIAGRAMS OF OLD ADDING MACHINES UNDERSTANDING DIGITAL ELECTRONICS, BY D. McWhorter. BOOLEAN CIRCUITS

UNDERSTANDING DIGITAL COMPUTERS. BY P. MIMS. A PERSONAL FAVORITE, BUT LOOK OUT FOR MISPRINTS!

INTRODUCTION TO MICROCOMPITERS, BY A. OSBORNE (4 VOLUMES). VERY DETAILED

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CHARLES BABBAGE, FATHER OF THE COMPUTER, BY CO. NALACY. 'AN EASY READ

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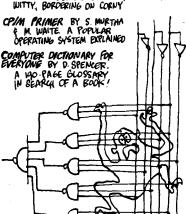


ILLUSTRATING BASIC BY D. ALCOCK. A QUICK COURSE, USING QUASI-CARTOONS

USING BASIC, BY R.DIDDAY & R.PAGE. A GENTILER, BUT WORDIER, APPROACH

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